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# AN ALGORITHM FOR OPTIMIZATION PROBLEMS

# WITH FUNCTIONAL INEQUALITY CONSTRAINTS

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### An Algorithm for Optimization Problems

## with Functional Inequality Constraints

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### Abstract

This paper presents an algorithm for minimizing an objective function subject to conventional inequality constraints as well as to inequality constraints of the functional type:  $\max_{\omega \in \Omega} \phi(z,\omega) \leq 0$ , where  $\Omega$  is a closed interval in R, and  $z \in \mathbb{R}^n$  is the parameter vector to be optimized. The algorithm is motivated by a standard earthquake engineering problem and the problem of designing linear multivariable systems. The stability condition (Nyquist criterion) and disturbance suppression condition for such systems are easily expressed as a functional inequality constraints.

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### 1. Introduction.

This paper presents an algorithm for solving problems of the form:

 $\min\{f^{o}(z) | f^{j}(z) \leq 0, j = 1, \dots, m; g^{j}(z) \leq 0, j = 1 \dots, r\}$ where  $z \in \mathbb{R}^{n}$  is the parameter vector to be optimized,  $f^{o}: \mathbb{R}^{n} \rightarrow \mathbb{R}$  is the criterion,  $g^{j}: \mathbb{R}^{n} \rightarrow \mathbb{R}$ ,  $j = 1, \dots, p$ , are conventional inequality constraints, and  $f^{j}: \mathbb{R}^{n} \rightarrow \mathbb{R}$ ,  $j = 1, \dots, m$ , are inequality constraints of the <u>functional</u> type, i.e.

$$f^{j}(z) \stackrel{\Delta}{=} \max_{\omega \in \Omega} \phi^{j}(z,\omega), \ j = 1 \dots m$$
(1)

where  $\Omega$  is a closed interval in R. The algorithm is motivated by problems arising in the design of earthquake resistant buildings [10] and in designing controllers for linear multivariable systems using frequency response techniques. For the latter, a sufficient condition for asymptotic stability of the closed-loop system is that (if the open-loop system is asymptotically stable) for all  $\omega \in \mathbb{R}$ , det(T(i $\omega$ ,z)) lies outside of a suitably specified subset B<sub>1</sub> of the complex plane (see Fig. 1), where det(T( $\cdot$ ,z)) is the determinant of the matrix return difference of the multivariable system corresponding to parameter z. If B<sub>1</sub> can be defined by:

$$B_{1} \stackrel{\Delta}{=} \{ s \in \bigcap | \theta(s) \leq 0 \}$$
(2)

where  $\theta: \mathbb{C} \rightarrow \mathbb{R}$ , then the stability functional constraint is:

$$f^{1}(z) \leq 0$$
  
e  $f^{1}(z) \stackrel{\triangle}{=} \max_{\omega \in \mathbb{R}} \phi^{1}(z,\omega), \text{ and } \phi^{1}(z,\omega) \stackrel{\triangle}{=} \theta(|T(i\omega,z)|)$ 

where

In practice it is known that there exists a  $\omega_c < \infty$  such that  $\phi^1(z,\omega) < 0$  for all feasible z, for all  $\omega \ge \omega_c$ , so that R may be replaced by  $\Omega = [0, \omega_c]$  in the definition of  $f^1$ .

Similarly, a <u>performance</u> constraint can be specified, for some  $\delta > 0$ , as  $\|T^{-1}(i\omega,z)\| \leq \delta$  for all  $\omega \in \Omega$ . Existing procedures for designing linear multivariable systems, when the state is not completely accessible, are either of the time-domain [1,2] or the frequency-domain type [3,4,5]. The former require relatively lengthy, complex calculations, and, usually, the specification of an artificial criterion (which may, itself, have to be repeatedly adjusted to meet basic design objectives); the latter require, at least in their current stage of development, considerable skill, and they do not necessarily yield the simplest designs. Since the latter procedures easily handle design objectives (e.g. stability and performance) which are meaningful in many practical situations, it seemed useful to develop an algorithm to achieve these objectives automatically. Standard algorithms cannot be employed since at least some of the constraints are of the functional type.

Because the problem contains functional constraints which cannot be evaluated without resorting to approximations, the algorithm is presented in two versions: the conceptual version, which assumes that the functionals can be evaluated exactly, and the implementable, which incorporates efficient approximation procedures compatible with convergence. The conceptual version is given primarily because it considerably simplifies the presentation.

In constructing these new algorithms we have made use of ideas appearing in Zontendijk's feasible directions algorithms [6] and in Demyanov's algorithm [7] for min max problems.

2. Definitions and Assumptions.

The problem considered is:

P1.  $\min\{f^{o}(z) | f^{j}(z) \leq 0, j \in J_{m}; g^{j}(z) \leq 0, j \in J_{p}\}$ where  $J_{k} \stackrel{\Delta}{=} \{1, 2, ..., k\}, k = m, p, f^{j}, j \in J_{m}$ , is defined in (1), and the following conditions are assumed to be satisfied: A1.  $f^{o}: R^{n} \rightarrow R, g^{j}: R^{n} \rightarrow R, j = 1 ... r$  are continuously differentiable. A2.  $\phi^{j}: \mathbb{R}^{n} \times \Omega \to \mathbb{R}, j = 1 \dots r$  are continuously differentiable. Let  $F \stackrel{\Delta}{=} \{z \in \mathbb{R}^{n} | f^{j}(z) \leq 0, j \in J_{m}; g^{j}(z) \leq 0, j \in J_{p}\}$  denote the <u>feasible set</u>, and  $\stackrel{F}{F}$  its interior. The algorithm is an extension of a method of feasible directions, and requires the determination of the " $\varepsilon$ -active constraints". Thus, let:

$$J_{\varepsilon}^{f}(z) \stackrel{\Delta}{=} \{j \in J_{m} | f^{j}(z) - \psi_{o}(z) \geq -\varepsilon\}$$
(3)

$$J_{\varepsilon}^{g}(z) \stackrel{\Delta}{=} \{j \in J_{p} | g^{j}(z) - \psi_{o}(z) \geq -\varepsilon\}$$
(4)

where  $\psi_{O}$ :  $\mathbb{R}^{n} \rightarrow \mathbb{R}$  is defined by:

$$\psi_{o}(z) \stackrel{\triangle}{=} \max\{0, \psi(z)\}$$
(5)

and  $\psi: \mathbb{R}^n \to \mathbb{R}$  is defined by:

$$\psi(z) \stackrel{\scriptscriptstyle {\stackrel{\circ}{=}}}{=} \max\{f^{j}(z), j \in J_{m}; g^{j}(z), j \in J_{p}\}$$
(6)

Clearly, if  $z \in F$ , then  $\psi_0(z) = 0$ . If  $j \in J_{\varepsilon}^{f}(z)$ , then  $\phi^{j}(z,\omega) - \psi_0(z) \geq -\varepsilon$ for some  $\omega \in \Omega$ ; the set of all such  $\omega$  will be denoted by:

$$\Omega_{\varepsilon}^{\mathbf{j}}(\mathbf{z}) \stackrel{\Delta}{=} \{ \omega \in \Omega | \phi^{\mathbf{j}}(\mathbf{z}, \omega) - \psi_{o}(\mathbf{z}) \geq -\varepsilon \}$$
(7)

The following further assumptions are made:

# A3. $\mathring{F} \neq \phi$ , $\ddot{\mathring{F}} = F$ .

A4.  $\Psi_z \in \mathbb{R}^n$ ,  $\Psi_j \in J_0^f(z)$ ,  $\Omega_0^j(z)$  is a finite set  $(\Omega_0^j(z) = \{\omega \in \Omega | \phi^j(z, \omega) = \Psi_0(z)\})$ A5.  $\Psi_z \in \mathbb{R}^n$ ,  $\Psi_\varepsilon > 0$ ,  $\Psi_j \in J_\varepsilon^f(z)$ ,  $\Omega_\varepsilon^j(z)$  is the union of a finite number of disjoint intervals  $g_{\varepsilon,k}^j(z)$ ,  $k = 1, \ldots k_\varepsilon^j(z)$ , possibly of zero length, i.e.

$$\Omega_{\varepsilon}^{j}(z) = \bigcup_{k \in \widetilde{C} \atop \varepsilon} I_{\varepsilon,k}^{j}(z) \qquad I_{\varepsilon,k}^{j}(z) \qquad (8)$$

(9)

where:  $\mathcal{K}_{\varepsilon}^{j}(z) \stackrel{\Delta}{=} \{1, 2, \ldots, \kappa_{\varepsilon}^{j}(z)\}$ 

The algorithm also uses a set of discrete frequencies  $\overline{\Omega}_{\varepsilon}^{j}(z) \subset \Omega_{\varepsilon}^{j}(z)$ defined  $\forall z \in \mathbb{R}^{n}, \forall j \in J_{\varepsilon}^{f}(z), by:$  $\overline{\Omega}_{\varepsilon}^{j}(z) \stackrel{\Delta}{=} \{\omega_{\varepsilon,k}^{j}(z) | k \in \mathbb{K}_{\varepsilon}^{j}(z)\}$  (10)

<sup>†</sup>The overbar denotes closure.

where  $\omega_{\varepsilon}^{j}k$  is the midpoint of  $I_{\varepsilon,k}^{j}$ ,  $\forall k \in K_{\varepsilon}^{j}(z)$ .

Most algorithms require directional derivatives of the constraint functions. This causes no difficulties for the conventional constraints  $g^{j}$ ,  $j \in J_{p}$  - the directional derivative corresponding to direction  $h \in \mathbb{R}^{n}$ is  $\langle \nabla g^{j}(z), h \rangle$ . To cope with the functional constraints, the function  $Df_{c}^{j}$ :  $\mathbb{R}^{n} \times \mathbb{R}^{n} \to \mathbb{R}$ , defined by:

$$Df_{\varepsilon}^{j}(z,h) \triangleq \max\{ \langle \nabla_{z} \phi^{j}(z,\omega), h \rangle, \omega \in \Omega_{\varepsilon}^{j}(z) \}, j \in J_{\varepsilon}^{f}(z)$$
$$\triangleq -\infty, j \notin J_{\varepsilon}^{f}(z)$$
(11)

 $\forall \epsilon \geq 0$ , is introduced. Similarly, for convenience,  $Dg_{\epsilon}^{j}: \mathbb{R}^{n} \times \mathbb{R}^{n} \rightarrow \mathbb{R}$  is defined by:

$$Dg_{\varepsilon}^{j}(z,h) \stackrel{\Delta}{=} \langle \nabla g_{j}(z),h \rangle, \ j \in J_{\varepsilon}^{g}(z)$$

$$\stackrel{\Delta}{=} -\infty, \ j \notin J_{\varepsilon}^{g}(z)$$
(12)

(13)

In choosing a search direction, a linear programming problem has to be solved. To make this low dimensional, the following 'approximation'  $\overline{Df}_{\varepsilon}^{j}: \mathbb{R}^{n} \times \mathbb{R}^{n} \to \mathbb{R}$  to  $Df_{\varepsilon}^{j}$  is employed:  $\overline{Df}_{\varepsilon}^{j}(z) \stackrel{\Delta}{=} \max\{ \langle \nabla_{z} \phi^{j}(z, \omega), h \rangle, \omega \varepsilon \widetilde{\Omega}_{\varepsilon}^{j}(z) \}, j \in J_{\varepsilon}^{f}(z)$ 

Because of A4, A5,  $Df_0^j = \overline{D}f_0^j$ ,  $\forall j \in J_m$ .

 $\stackrel{\Delta}{=} -\infty$ , j  $\notin J_{2}^{f}(z)$ 

For given  $z \in R^n$ ,  $\varepsilon > 0$ , the search direction  $h_{\varepsilon}(z) \in R^n$  is obtained by solving the following linear programming sub problem:

P2. Compute  $h_{\epsilon}(z) \in R^{n}$  by solving:

$$\overline{\theta}_{\varepsilon}(z) = \min_{h \in S} \max\{-\psi_{o}(z) + \langle \nabla f^{o}(z), h \rangle; \\
\overline{D}f_{\varepsilon}^{j}(z, h), j \in J_{m}; Dg_{\varepsilon}^{j}(z, h), j \in J_{p}\}$$
(14)

where:

$$S \stackrel{\Delta}{=} \{h \in R^{L} \mid \|h\|_{\infty} \leq 1\}$$
(15)

The function  $\psi_0$  makes the algorithm concentrate initially on driving z into the feasible region F. In the conceptual algorithm, once z becomes

feasible it remains feasible.

To check that the computed search direction  $h_\epsilon(z)$  is satisfactory,  $\tilde{\theta}_\epsilon(z),$  defined by:

$$\tilde{\theta}_{\varepsilon}(z) = \max\{-\psi_{o}(z) + \langle \nabla f^{o}(z), h_{\varepsilon}(z) \rangle; \\Df_{\varepsilon}^{j}(z,h_{\varepsilon}(z)), j \in J_{m}; Dg_{\varepsilon}^{j}(z,h_{\varepsilon}(z)), j \in J_{p}\}$$
(16)

is evaluated. To aid the analysis of the algorithm, the function  $\theta_{\epsilon}$ :  $\mathbb{R}^n \to \mathbb{R}$ , defined,  $\forall \epsilon \geq 0$ , by:

$$\theta_{\varepsilon}(z) = \min_{h \in S} \max\{-\psi_{o}(z) + \langle \nabla f^{o}(z), h \rangle; Df_{\varepsilon}^{j}(z, h), j \in J_{m}; \\ h \in S ; Dg_{\varepsilon}^{j}(z, h), j \in J_{p}\}$$
(17)

is introduced.  $e_0(z) = 0$  is a necessary condition of optimality.

The final assumption can now be stated:

A6. 
$$\Psi_z \in \mathbb{R}^n$$
,  $\{\nabla_z \phi^j(z, \omega), \omega \in \Omega_0^j(z), j \in J_0^f(z); \nabla g^j(z), j \in J_0^g(z)\}$  is a set of linearly independent vectors.<sup>†</sup>

This assumption is not easily verifiable. However, it indicates that the algorithm may jam up, even when  $z \in F$ , at a point where it is not possible to reduce all the active constraints. This might indicate that the problem is not well posed, and that some constraints should be removed. The assumption ensures that there exists an  $h \in \mathbb{R}^n$  such that  $Df_0^j(z,h) < 0$  $\forall j \in J_m$  and  $Dg_0^j(z,h) < 0$   $\forall j \in J_p$ ,  $\forall z \in F^c$ .

3. The Conceptual Algorithm.

The conceptual algorithm can now be stated:

Algorithm I:

<sup>&</sup>lt;sup>†</sup>This assumption is related to the Kuhn-Tucker constraint qualification.

<u>Step 3</u>: If  $\overline{\theta}_{\varepsilon}(z_i) \leq -2\varepsilon$ , go to Step 5.

Else, set  $\varepsilon = \varepsilon/2$  and go to Step 4.

<u>Step 4</u>: If  $\varepsilon \leq \varepsilon_1$  and  $\theta_0(z_i) = 0$ , stop. Else go to Step 2.<sup>+</sup>

<u>Step 5</u>: If  $\tilde{\theta}_{\varepsilon}(z_i) \leq -\varepsilon$  go to Step 5. Else set  $\varepsilon = \varepsilon/2$  and go to Step 2.

Step 6: If 
$$z_i \in F$$
 compute the smallest integer  $l_i$  such that:

$$f^{0}(z_{i} + \beta h_{\varepsilon}(z_{i})) - f^{0}(z_{i}) \leq -\beta \alpha \varepsilon$$
(18)

$$f^{j}(z_{i} + \beta^{z_{i}}h_{\varepsilon}(z_{i})) \leq 0, \quad \forall j \in J_{m}$$
(19)

$$g^{j}(z_{i} + \beta^{k_{i}}h_{\varepsilon}(z_{i})) \leq 0, \quad \forall j \in J_{p}$$
 (20)

If  $z_i \in F^c$  (the complement of F in  $R^n$ ) compute the smallest

integer  $l_i$  such that:

$$\psi(z_{i} + \beta^{l_{i}}h_{\varepsilon}(z_{i})) - \psi(z_{i}) \leq -\beta^{l_{i}}\alpha\varepsilon$$
(21)

Step 7: Set 
$$z_{i+1} = z_i + \beta^{i}h_{\epsilon}(z_i)$$
,  $i = i + 1$ . Go to Step 1 (Step 2).<sup>††</sup>

For future reference let  $\overline{\varepsilon}(z)$  denote the value of  $\varepsilon$  generated by the algorithm (with "go to Step 1" in Step 7) given  $z_i = z$ , i.e.  $\forall z \in \mathbb{R}^n$ ,  $\overline{\varepsilon}(z)$  is the largest number, of the form  $\varepsilon_0/2^k$ , where k is a non-negative integer, satisfying:

$$\frac{\overline{\theta}_{\varepsilon}(z)}{\widetilde{\theta}_{\varepsilon}(z)}(z) \leq -2 \overline{\varepsilon}(z)$$
(22)
$$\frac{\widetilde{\theta}_{\varepsilon}(z)}{\widetilde{\theta}_{\varepsilon}(z)}(z) \leq -\overline{\varepsilon}(z)$$
(23)

Note that Steps 1-7 of the conceptual algorithm (with "go to Step 1" in Step 7) define a Map A:  $\mathbb{R}^n \to 2^{\mathbb{R}^n}$ , so that  $z_{i+1} \in A(z_i)$ , with the property that  $A(F) \subset F$ .

<sup>†</sup>Note that if  $\varepsilon > \varepsilon_1$ , there is no need to evaluate  $\theta_0(z_i)$ .

<sup>&</sup>lt;sup>††</sup>We shall give a proof of convergence for the case "go to Step 1" in Step 7, i.e., when  $\varepsilon$  is reset to  $\varepsilon_0$  at each iteration. The algorithm also converges when  $\varepsilon$  is forced to decrease monotonically, i.e. "go to Step 2" is used in Step 7. However, the proof of convergence is then substantially more complicated and will therefore be omitted.

The algorithm essentially has two different cost functions, one for F and one for  $F^{C}$ . In order to analyze its convergence properties we make use of the following general convergence theorem for algorithm models with this structure. (Theorem 1 below generalizes theorem 1.3.10 in [8]). Theorem 1.

Suppose an algorithm A:  $\mathbb{R}^n \to 2^{\mathbb{R}^n}$ , with associated cost function c:  $\mathbb{R}^n \to \mathbb{R}$ , has the following properties (with  $\mathbb{N}_{\varepsilon}(z) \triangleq \{z' | \|z-z'\| \le \varepsilon\}$ ): 1)  $\exists F \subset \mathbb{R}^n$  such that  $c = c^1$  on F and  $c = c^2$  on  $F^c$ , where  $c^1$ ,  $c^2$ :  $\mathbb{R}^n \to \mathbb{R}$ are continuous.

2)  $A(F) \subseteq F$ .

3)  $\forall z \in \mathbb{R}^{n}$ , such that z is not desirable,  $\exists \varepsilon(z) > 0$ ,  $\exists \delta(z) < 0$  such that (i)  $c^{1}(z'') - c^{1}(z') \leq \delta(z) < 0$ ,  $\forall z' \in \mathbb{N}_{\varepsilon(z)}(z) \cap F$ ,  $\forall z'' \in \mathbb{A}(z')$ (ii)  $c^{2}(z'') - c^{2}(z') \leq \delta(z) < 0$ ,  $\forall z' \in \mathbb{N}_{\varepsilon(z)}(z) \cap F^{c}$ ,  $\forall z'' \in \mathbb{A}(z')$ 

Then, every accumulation point of an infinite sequence  $\{z_i\}$  generated according to the rule  $z_{i+1} = y_i \in A(z_i)$  if  $c(y_i) < c(z_i)$ ,  $z_{i+1} = z_i$  otherwise, i = 0, 1, 2, ..., is desirable.<sup>†</sup>

<u>Proof</u>. Suppose  $\{z_{\underline{i}}\}$  is an infinite sequence, generated as above, which has a subsequence, indexed by K, converging to a  $z \notin \Delta$ , where  $\Delta$  denotes the set of desirable points. From assumption 2), either there exists a finite k such that  $z_{\underline{i}} \in F \forall \underline{i} \geq k$ , or  $z_{\underline{i}} \in F^{C} \forall \underline{i}$ . In the first case  $(\underline{j} = 1) c^{1}(z_{\underline{i}}) \stackrel{K}{\rightarrow} c^{1}(z)$  (m in the second  $(\underline{j} = 2) c^{2}(z_{\underline{i}}) \stackrel{K}{\rightarrow} c^{2}(z)$ . Hence,  $\exists k' \in [0, \infty)$  such that, for  $\underline{j} = 1$  or 2:

 $c^{j}(z_{i+1}) - c^{j}(z_{i}) \leq \delta(z), \quad \forall i \in K \text{ such that } i \geq k'.$ Hence for j = 1 or 2, (assuming, without loss of generality, that  $z_{0} \in F$ if  $z_{i} \in F$  for all i sufficiently large)  $c^{j}(z) = c^{j}(z_{i}) = \sum_{i=1}^{n} [c^{j}(z_{i}) - c^{j}(z_{i})]$ 

$$c^{J}(z) - c^{J}(z_{0}) = \sum_{\substack{i \in K \\ i \geq k'}} [c^{J}(z_{i+1}) - c^{J}(z_{i})]$$

<sup>†</sup>Note that if the sequence is finite, i.e.  $z_i = z_k$  for all  $i \ge l$ , then  $z_k$  is desirable.

+ 
$$\sum_{\substack{i \in K \\ i < k'}} [c^{j}(z_{i+1}) - c^{j}(z_{i})]$$
  
+ 
$$\sum_{i \notin K} [c^{j}(z_{i+1}) - c^{j}(z_{i})]$$

since each term in the first summation is less than or equal to  $-\delta(z)$  and all the other terms are negative. This contradicts the convergence of  $\{c^{j}(z_{i})\}, j = 1 \text{ or } 2$ . Hence z is desirable. #

To establish that the conceptual algorithm satisfies the assumptions of Theorem 1, we require the following results, which are proven in the Appendix.

<u>Proposition 1</u>. If  $z \in F$  is optimal for P1, then  $\theta_0(z) = \overline{\theta}_0(z) = \widetilde{\theta}(z) = 0$ . # <u>Proposition 2</u>.  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\neg \gamma > 0$ ,  $\neg \varepsilon > 0$  such that:  $\overline{\varepsilon}(z') \ge \varepsilon$ ,  $\forall z' \in \mathbb{N}_{\gamma}(z)$ . #

<u>Proposition 3</u>.  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0, -] \gamma > 0, -] \varepsilon > 0$  such that:

 $f^{o}(z') - f^{o}(z') \leq -\varepsilon, \quad \forall z' \in N_{\gamma}(z) \cap F, \quad \forall z'' \in A(z')$ 

$$\begin{split} \psi(z'') - \psi(z') &\leq -\varepsilon, \quad \forall z' \in \mathbb{N}_{\gamma}(z) \cap F^{C}, \quad \forall z'' \in (A(z')), \\ \text{where A: } \mathbb{R}^{n} \rightarrow 2^{\mathbb{R}^{n}} \text{ is defined by Steps 1 to 7 of the conceptual algorithm,} \\ \text{with "go to Step 1" in Step 7.} \qquad \qquad \# \end{split}$$

As a consequence of Proposition 1 we define the set  $\Delta$  of desirable points as follows:

$$\Delta \stackrel{\Delta}{=} \{ z \in \mathbb{R}^n | \theta_0(z) = 0 \}$$
(24)

Since algorithm I stops only at  $z_i \in \Delta$  we only need to consider the case where the sequence  $\{z_i\}$  is infinite. Our first main result is: Theorem 2.

Consider algorithm I with "go to Step 1" in Step 7. Then every accumulation point  $z^*$  of an infinite sequence  $\{z_i\}$  generated by algorithm I is

desirable.

<u>Proof.</u> Since  $f^{\circ}$ :  $\mathbb{R}^{n} \to \mathbb{R}$  and  $\psi$ :  $\mathbb{R}^{n} \to \mathbb{R}$  are continuous, the conceptual algorithm satisfies condition 1) of Theorem 1 will  $c^{1} \triangleq f^{\circ}$ ,  $c^{2} \triangleq \psi$ .

It is obvious from Step 5 of the conceptual algorithm that if  $z_i \in F$ , then  $z_{i+1} \in F$  and consequently, defining A:  $\mathbb{R}^n \to 2^{\mathbb{R}^n}$  by Steps 1-7, we find that  $A(z) \subset F$ . Condition 2) of Theorem 1 is therefore satisfied.

Finally, Proposition 3 shows that the conceptual algorithm satisfies condition 3) of Theorem 1. #

We state without proof a similar result for the alternative version of the conceptual algorithm.

### Theorem 3.

Consider algorithm I with "go to Step 2" in Step 7. Then every accumulation point  $z^*$  of an infinite sequence  $\{z^i\}$  generated by algorithm I is desirable.

The interested reader can prove Theorem 3, by modifying a convergence theorem due to Klessig [9], analogously to the way we have modified Theorem 1.3.10. in [8].

## 4. The Implementable Algorithm.

The algorithm below takes into account the fact that it is impossible to evaluate expressions such as max  $\phi^j(z,\omega)$ , with high precision, a large  $\omega \in \Omega$ number of times, without incurring high computational costs. We assume that  $\Omega$  is a closed interval,

$$\Omega = [\omega_0, \omega_c]$$
(25)

and, for  $q = 1, 2, \ldots, \epsilon \ge 0$ , we define

$$\Omega_{q} \stackrel{\Delta}{=} \{ \omega_{\ell} | \omega_{\ell} = \omega_{0} + \ell [\omega_{c} - \omega_{0}]/2^{q}, \ \ell = 0, 1, ..., 2^{q} \}$$
(26)

$$f_{q}^{j}(z) \stackrel{\Delta}{=} \max_{\omega \in \Omega} \phi^{j}(z,\omega), j = 1, 2, ..., m$$
(27)

$$\mathbf{F}_{q} \stackrel{\Delta}{=} \{ \mathbf{z} \in \mathbb{R}^{n} | \mathbf{f}_{q}^{j}(\mathbf{z}) \leq 0, \ \mathbf{j} \in \mathbf{J}_{m}^{j}; \ \mathbf{g}^{j}(\mathbf{z}) \leq 0, \ \mathbf{j} \in \mathbf{J}_{p}^{j} \}$$
(28)

$$\psi_{q}(z) \stackrel{\Delta}{=} \max\{f_{q}^{j}(z), j \in J_{m}, q^{j}(z), j \in J_{p}\}$$
(29)

$$\psi_{q,o}(z) \stackrel{\Delta}{=} \max\{0, \psi_{q}(z)\}$$
(30)

$$J_{q,\varepsilon}^{f}(z) \stackrel{\Delta}{=} \{j \in J_{m} | f_{q}^{j}(z) - \psi_{q,o}(z) \geq -\varepsilon\}$$
(31)

$$J_{q,\varepsilon}^{g}(z) \stackrel{\Delta}{=} \{j \in J_{p} | g^{j}(z) - \psi_{q,o}(z) \geq -\varepsilon\}$$
(32)

$$\Omega_{q,\varepsilon}^{j}(z) = \{ \omega \in \Omega_{q} | \phi^{j}(z,\omega) - \psi_{q,o}(z) \ge -\varepsilon \}, \ j \in J_{m},$$
(33)

$$\overline{\Omega}_{q,\varepsilon}^{j}(z) = \{ \omega_{q\varepsilon,k}^{j}(z) | k \in \widetilde{\mathcal{K}}_{q,\varepsilon}^{j}(z) \}$$
(34)

where the  $\omega_{q,\varepsilon,k}^{j}(z)$  are mid points of intervals  $I_{q,\varepsilon,k}^{j}(z) = [\omega',\omega'']$ , where  $\omega', \omega'' \in \Omega_{q,\varepsilon}^{j}(z), \omega' = \omega''$  is allowed, and for any  $\omega^{*} \in \Omega_{q,\varepsilon}^{j}(z), \omega^{*} \notin I_{q,\varepsilon,k}^{j}(z)$ min  $\omega \in I_{q,\varepsilon,k}^{j}(z)$  $|\omega - \omega^{*}| \ge 2 \cdot [(\omega_{c} - \omega_{o})/2^{q}].$ 

Let  $I^{\dagger} \triangleq \{0, 1, 2, ...\}$ , then  $\forall j \in J_m$ ,  $\forall q \in I^{\dagger}$ ,  $\forall \epsilon \ge 0$ ,  $Df_{q,\epsilon}^{j}$  is defined as in (11) with  $\Omega_{q,\epsilon}^{j}$  replacing  $\Omega_{\epsilon}^{j}$ , and  $\overline{D}f_{q,\epsilon}^{j}$  is defined as in (13) with  $\overline{\Omega}_{q,\epsilon}^{j}$  replacing  $\overline{\Omega}_{\epsilon}^{j}$ . Next,  $\forall q \in I^{\dagger}$ ,  $\forall \epsilon \ge 0$ ,  $\forall z \in \mathbb{R}^{n}$ ,  $h_{q,\epsilon}(z)$  and  $\overline{\theta}_{q,\epsilon}(z)$ are obtain by solving:

P3. 
$$\overline{\theta}_{q,\varepsilon}(z) = \min_{h \in s} \max\{-\psi_{q,o}(z) + \langle \nabla f^{o}(z), h \rangle;$$
  
 $\overline{D}f_{q,\varepsilon}^{j}(z,h), j \in J_{m}; Dg_{\varepsilon}^{j}(z,h), j \in J_{p}\}$  (35)

Finally,  $\forall q \in I^+$ ,  $\forall \epsilon \geq 0$ ,  $\tilde{\theta}_{q,\epsilon}$ :  $\mathbb{R}^n \rightarrow \mathbb{R}$  is defined as in (16) with  $Df_{q,\epsilon}^j$ replacing  $Df_{\epsilon}^j$ ,  $\forall j \in J_m$ .

Algorithm II.

<u>Data</u>:  $z_0 \in \mathbb{R}^n$ ,  $\alpha \in (0,1)$ ,  $\beta \in (0,1)$ ,  $\varepsilon_0 \in (0,\infty)$   $0 < \varepsilon_1 << \varepsilon_0$ ,  $q_0 > 0$ , an integer.

Step 0: Set i = 0,  $q = q_0$ .

# <u>Step 1</u>: Set $\varepsilon = \varepsilon_0$ .

- <u>Step 2</u>: Compute  $h_{q,\varepsilon}(z_i)$  and  $\overline{\theta}_{q,\varepsilon}(z_i)$  by solving P3.
- <u>Step 3</u>: If  $\overline{\theta}_{q,\epsilon}(z_i) \leq -2\epsilon$ , go to Step 4. Else set  $\epsilon = \epsilon/2$  and go to Step 2. <u>Step 4</u>: If  $\epsilon \leq \epsilon_1$  and  $\overline{\theta}_{q,0}(z_i) = 0$ , set  $y_q = z_i$ , set q = q + 1 and go to Step 1. Else go to Step 5.

<u>Step 5</u>: If  $\tilde{\theta}_{q,\epsilon}(z_i) \leq -\epsilon$ , go to Step 6. Else set  $\epsilon = \epsilon/2$  and go to Step 2. <u>Step 6</u>: If  $z_i \in F_a$  compute the smallest integer  $\ell_i$  such that

$$f^{o}(z_{i} + \beta^{\ell} h_{q,\epsilon}(z_{i})) - f^{o}(z_{i}) \leq -\beta^{\ell} \alpha \epsilon$$
(36)

$$f_{q}^{j}(z_{i} + \beta_{q,\varepsilon}^{\prime i}(z_{i})) \leq 0 \quad \forall j \in J_{m}$$
(37)

$$g^{j}(z_{i} + \beta^{i}h_{q,\varepsilon}(z_{i})) \leq 0 \quad \forall j \in J_{p}$$
 (38)

If  $z_i \in F_q^c$  (the complement of  $F_q$  in  $R^n$ ) compute the smallest integer  $l_i$  such that

$$\psi_{q}(z_{i} + \beta^{l}h_{q,\varepsilon}(z_{i})) - \psi_{q}(z_{i}) \leq -\beta^{l}i_{\alpha\varepsilon}$$
(39)

<u>Step 7</u>: Set  $z_{i+1} = z_i + \beta^{-1}h_{q,\epsilon}(z_i)$ , i = i+1. <u>Step 8</u>: If  $z_i \in F_q$  and  $\epsilon \leq \epsilon_1$ , go to Step 9. Else go to Step 2. <u>Step 9</u>: If  $\overline{\theta}_{q,0}(z_i) \geq -1/2^q$ , set  $y_q = z_i$ , set q = q + 1 and go to Step 1. Else go to Step 2.

The following result follows readily from the results in the appendix and the assumptions made in the beginning of the paper.

### Theorem 4.

Every accumulation point  $y^*$  of the sequence  $\{y_q\}$  constructed by algorithm II satisfies  $y^* \in F$  and  $\theta_0(y^*) = 0$ . #

## 5. <u>Conclusion</u>.

It is hoped that the algorithms presented in this paper will be of use for problems with functional inequality constraints. One example of this class of problems was mentioned in the introduction: the design of

linear multivariable systems when stability and performance constraints are expressed in terms of frequency response functionals. Another example is the design of buildings is withstand earthquakes [10]; here the cost would be the weight of the building, and a typical constraint the maximum deviation of a floor during an earthquake.

For the problem of designing linear multivariable systems, a set  $C_1$  has to be chosen so that det(T(i $\omega$ ))  $\notin B_1 \quad \forall \ \omega \in \Omega$  is sufficient to ensure "relative", as well as asymptotic, stability i.e. the locus det(T(i $\Omega$ )) (det(T(i $\Omega$ ))  $\triangleq$  { det(T(s))s = i $\omega$ ,  $\omega \in \Omega$ }) should maintain a reasonable distance from the origin of the complex plane . It is probably worth mentioning that a prescribed degree of "relative stability" may be achieved by considering rather the locus det(T(( $\alpha + i$ ) $\Omega$ )), where  $\alpha < 0$  is suitably chosen, in place of det(T(i $\Omega$ )); if this is done  $C_1$  may be chosen merely to prevent encirclement by det(T(( $\alpha + i$ ) $\Omega$ )) of the origin.

The algorithm has a feature which should be of interest also for conventional optimization problems (m = 0), i.e. without functional constraints, namely that the initial points  $z_0$  need not be feasible. Algorithm 1 automatically drives z into the feasible region and then maintains feasibility; i.e. unlike conventional feasible directions algorithms, it solves the "phase I" and "phase II" problems simultaneously.

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### Appendix

<u>Proposition 1</u>. If  $z \in F$  is optimal for P1, then  $\theta_0(z) = \overline{\theta}_0(z) = \tilde{\theta}(z) = 0$ . <u>Proof.</u> (i) It follows from assumption A4, (7), (10) that,  $\forall z \in \mathbb{R}^m$ ,  $\forall j \in J_m$ ,  $\Omega_0^j(z) = \overline{\Omega}_0^j(z)$ , and hence from (11), (13), that  $\forall j \in J_m$ ,  $Df_0^j = \overline{Df}_0^j$ . It now follows from (14), (16), (17) that  $\forall z \in \mathbb{R}^n$ ,  $\overline{\theta}_0(z) = \tilde{\theta}_0(z) = \theta_0(z)$ . (ii) If  $z \in F$  is optimal, and  $\theta_0(z) \leq \delta < 0$ , then  $\psi_0(z) = 0$ , and  $\exists h \in S$  such that:

$$\begin{aligned} \mathrm{Df}_{0}^{\mathbf{j}}(\mathbf{z},\mathbf{h}) &\leq \delta \quad \forall \mathbf{j} \in \mathbf{J}_{\mathbf{m}} \\ \mathrm{Dg}_{0}^{\mathbf{j}}(\mathbf{z},\mathbf{h}) &\leq \delta \quad \forall \mathbf{j} \in \mathbf{J}_{\mathbf{p}}. \end{aligned}$$

It follows from the continuity of  $\nabla f^{0}$ ,  $\nabla_{z}\phi^{j}$ ,  $\forall j \in J_{m}$  and  $\nabla g^{j}$ ,  $\forall j \in J_{p}$ , that  $\neg \lambda > 0$  such that  $z + \lambda$  h  $\in$  F and  $f^{0}(z + \lambda h) - f^{0}(z) < 0$ , contradicting the optimality of z. #.

Fact 1. (i) 
$$\forall z \in \mathbb{R}^n$$
,  $\overline{\theta}_0(z) = \tilde{\theta}_0(z) = \theta_0(z)$ 

(ii) 
$$\forall z \in \mathbb{R}^n$$
,  $\forall \varepsilon \geq 0$ ,  $\overline{\theta}_{\varepsilon}(z) \leq \theta_{\varepsilon}(z) \leq \tilde{\theta}_{\varepsilon}(z)$ 

(iii) 
$$\forall z \in \mathbb{R}^n$$
, if  $\varepsilon_1 \leq \varepsilon_2$  then  $\theta_{\varepsilon_1}(z) \leq \theta_{\varepsilon_2}(z)$ .

<u>Proof</u>. (i) is proven in Proposition 1.

(ii) follows from the fact that  $\forall z \in \mathbb{R}^n$ ,  $\forall \varepsilon \geq 0$ ,  $\forall j \in J_m$ ,  $\overline{\Omega}_{\varepsilon}^j(z) \subset \Omega_{\varepsilon}^j(z)$ . (iii) follows from the fact that  $\varepsilon_1 \leq \varepsilon_2$  implies that

$$\Omega_{\varepsilon_1}^{\mathbf{j}}(\mathbf{z}) \subset \Omega_{\varepsilon_2}^{\mathbf{j}}(\mathbf{z}) \qquad \#$$

<u>Fact 2</u>.  $\forall z \in \mathbb{R}^n$ ,  $\forall j \notin J^f_{\varepsilon}(z) \ (\forall j \notin J^g_{\varepsilon}(z)) \exists \gamma > 0$  such that  $j \notin J^f_{\varepsilon}(z') \ (j \notin J^f_{\varepsilon}(z')) \ \forall z' \in \mathbb{N}_{\gamma}(z).$ 

<u>Proof.</u> If  $j \notin J_{\varepsilon}^{f}(z)$ , then  $\exists \delta > 0$  such that  $\delta^{j}(z) - \psi_{0}(z) \leq -\varepsilon - \delta$ .

The result follows directly from the continuity of  $f^{j}$ ,  $\psi_{o}$ . The proof for the second case is similar. ⋕. <u>Fact 3</u>.  $\forall z \in \mathbb{R}^n$ ,  $\forall \varepsilon \geq 0$ ,  $\forall j \in J_{\varepsilon}^{f}(z)$ ,  $\forall \delta > 0 = \gamma > 0$  such that  $\mathfrak{M}_{\varepsilon}^{j}(z') \subseteq \mathtt{N}_{\delta}(\mathfrak{\Omega}_{\varepsilon}^{j}(z)) \ \forall z' \in \mathtt{N}_{\gamma}(z), \ \text{where} \ \mathtt{N}_{\delta}, \ \mathtt{N}_{\gamma}, \ \text{denote, respectively, } \delta,$  $\gamma$  neighborhoods of  $\Omega_{F}^{j}(z)$ , and z. Proof. First, it follows from (5) and (6) that  $\psi_{\Omega}: \mathbb{R}^{n} \to \mathbb{R}$  is continuous. If  $\mathbb{N} \stackrel{\Delta}{=} \mathbb{N}_{\delta}(\Omega_{\varepsilon}^{j}(z)) = \Omega$ , the result is trivially true. If  $N \neq \Omega$ , then  $N^{C} \subset (\Omega_{\varepsilon}^{j}(z))^{C}$ , and  $\exists \eta > 0$  such that  $\forall j \in J_{\varepsilon}^{f}(z)$  sup  $\omega \in N^{C}$  $[\phi^{j}(z,\omega)-\psi_{o}(z)] \leq -\varepsilon - \eta. \quad (\text{for, if not, } \sup_{\omega \in N^{C}} [\phi^{j}(z,\omega)-\psi_{o}(z)] \geq -\varepsilon,$ and  $\exists$  a sequence  $\{\omega_{\ell}\}$  such that  $\omega_{\ell} \in \mathbb{N}^{C}$  and  $[\phi^{j}(z,\omega_{\ell}) - \psi_{O}(z)] < -\varepsilon$  $\forall \ell \geq 0 \text{ and } \omega_{\ell} \neq \omega^* \in \overline{N}^C \text{ where } \omega^* \text{ satisfies } [\phi^j(z,\omega^*) - \psi_o(z)] \geq -\varepsilon;$ however this implies that  $\omega^* \in \Omega_{\varepsilon}^j(z)$ , i.e.  $\Omega_{\varepsilon}^j(z) \cap \overline{N}^c \neq \phi$ , contradicting the fact that  $\Omega_{f}^{j}(z) \subset int N$ ). Hence  $\frac{1}{2} \gamma > 0$  such that:  $\sup_{\omega \in \mathbb{N}^{c}} [\phi^{j}(z',\omega) - \psi_{o}(z')] = \max_{\omega \in \overline{\mathbb{N}}^{c}} [\phi^{j}(z',\omega) - \psi_{o}(z')] \leq -\varepsilon - \eta/2 < -\varepsilon$  $\forall z' \in N_{\gamma}(z)$ . Thus  $\Omega_{\varepsilon}^{j}(z') \subseteq N$ ,  $\forall z' \in N_{\gamma}(z)$ . # <u>Fact 4</u>.  $\forall z \in \mathbb{R}^n$ ,  $\forall j \in J_0^f(z)$ ,  $\Omega_{\varepsilon}^j(z) \to \Omega_0^j(z)$ , with respect to inclusion, as  $\varepsilon \rightarrow 0$ . Proof. The result follows from assumption A4, and the continuity of  $\phi^{j} - \psi_{o}, \forall j \in J_{m}.$ # <u>Corrollary</u>.  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\exists \varepsilon > 0$  such that  $\theta_{\varepsilon'}(z) < 0$ , ¥ ε' ∈ [0,ε]. Fact 5.  $\forall z \in R^n$ ,  $\forall \delta > 0$ ,  $\forall \gamma > 0$ ,  $\boxed{} \epsilon > 0$  such that  $\mu(\mathfrak{u}_{\varepsilon}^{j},(\mathbf{z}')) \leq \delta, \quad \forall \mathbf{z}' \in \mathbb{N}_{\gamma}(\mathbf{z}), \quad \forall \mathbf{\varepsilon}' \in [0,\varepsilon], \quad \forall \mathbf{j} \in \mathbb{J}_{m},$ where  $\mu(\cdot)$  denotes measure (total length).

<u>Proof</u>. If  $j \notin J_{c}^{f}$ , (z'),  $\mu(\Omega_{c}^{j}, (z')) = 0$ . Let  $j \in J_{c}^{f}, (z')$  and assume that Fact 5 is false, i.e.  $\sup_{z' \in N_{\nu}(z)} \mu(\Omega_{\varepsilon}^{j},(z')) > \delta, \quad \forall \epsilon' > 0. \text{ Hence } \exists$ sequences  $\{z_i\}$ , with  $z_i \rightarrow z^*$ , and  $\{\varepsilon_i\}$ , with  $\varepsilon_i \rightarrow 0$ , such that  $z_i \in N_v(z)$ and  $\mu(\Omega_{\varepsilon_1}^{\prime}(z_1)) \geq \delta/2$ ,  $\forall i \geq 0$ . But, by Fact 4, -]  $\varepsilon^* > 0$  such that  $\mu(\Omega_{*}^{j}(z^{*})) \leq \delta/8$ , and  $\exists i^{*}$  such that  $\varepsilon_{i} < \varepsilon^{*} \forall i \geq i^{*}$ , and (by Fact 3):  $\Omega_{\varepsilon_{i}}^{j}(z_{i}) \subseteq \Omega_{\varepsilon}^{j}(z_{i}) \subseteq \mathbb{N}_{\delta/8}(\Omega_{\varepsilon}^{j}(z^{*}))$  $\forall i \geq i^*$ . Since  $\mu(\Omega_{\epsilon}^j(z^*)) \leq \delta/8$  we have  $\mu(\Omega_{\epsilon}^j(z_i)) \leq \delta/4$ ,  $\forall i \geq i^*$ . But this contradicts  $\mu(\Omega_{\varepsilon_1}(z_1)) \geq \delta/2, \quad \forall i \geq 0.$ ∦. Fact 6.  $\forall z \in \mathbb{R}^n$ ,  $\forall h \in S$ ,  $\forall \delta > 0$ ,  $\neg \gamma > 0$ ,  $\neg \varepsilon > 0$  such that:  $\mathrm{Df}_{\varepsilon}^{j},(z',h) \leq \overline{\mathrm{D}}\mathrm{f}_{\varepsilon}^{j},(z',h) + \delta, \forall z' \in \mathrm{N}_{\mathrm{v}}(z), \forall \varepsilon' \in [0,\varepsilon], \forall j \in \mathrm{J}_{\mathrm{m}}(\mathrm{A}1).$ <u>Proof</u>. If  $j \notin J_{\epsilon}^{f}(z')$ , then  $j \notin J_{\epsilon}^{f}(z')$ , for  $\epsilon' \in [0, \epsilon]$ , and  $Df_{c}^{j}(z',h) = -\infty$  and (Al) is trivially satisfied. From the continuity of  $\{\nabla_{\mathbf{z}}\phi^{\mathbf{j}}, \mathbf{j} \in J_{\mathbf{m}}\}, \quad \forall \mathbf{z} \in \mathbb{R}^{n}, \quad \forall \delta > 0, \quad \exists \delta_{1} > 0, \quad \exists \gamma_{1} > 0 \text{ such that:}$  $\|\nabla_{\mathbf{z}}\phi^{\mathbf{j}}(\mathbf{z}',\omega') - \nabla_{\mathbf{z}}\phi^{\mathbf{j}}(\mathbf{z},\omega)\| \leq \delta/\sqrt{n}$ 

<u>Proof</u>. The result follows from the definitions of  $\overline{\theta}_{\varepsilon}$ ,  $\tilde{\theta}_{\varepsilon}$ , the continuity of  $\psi_{o}$  and  $\nabla f^{o}$ , and Fact 6.

Fact 8. 
$$\forall z \in \mathbb{R}^n$$
 such that  $\theta_0(z) < 0$ ,  $\exists \gamma > 0$ ,  $\exists \varepsilon > 0$  such that:  
 $\theta_{\varepsilon'}(z') \leq \frac{1}{2} \theta_{\varepsilon}(z) < 0 \quad \forall z' \in \mathbb{N}_{\gamma}(z), \quad \forall \varepsilon' \in [0, \varepsilon]$ 

<u>Proof</u>. By Fact 4,  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\exists \varepsilon > 0$  such that  $\theta_{\varepsilon}(z) < 0 \quad \forall \varepsilon' \in [0, \varepsilon]$ . Now:

$$\begin{aligned} & \theta_{\varepsilon}(z) \stackrel{\Delta}{=} \min_{h \in S} \max\{-\psi_{o}(z) + \langle \nabla f^{o}(z), h \rangle; \langle \nabla_{z} \phi^{j}(z, \omega), h \rangle, \\ & h \in S \\ & \omega \in \Omega_{\varepsilon}^{j}(z), j \in J_{\varepsilon}^{f}(z); Dg_{\varepsilon}^{j}(z), j \in J_{p} \end{aligned}$$

Let  $\delta > 0$  be such that  $\hat{\theta}_{\varepsilon}(z) \leq 3/4 \ \theta_{\varepsilon}(z)$ , where  $\hat{\theta}_{\varepsilon} \colon \mathbb{R}^{n} \to \mathbb{R}$  is defined by:  $\hat{\theta}_{\varepsilon}(z') \stackrel{\Delta}{=} \min_{h \in S} \max\{-\psi_{0}(z') + \langle \nabla f^{0}(z'), h \rangle; \langle \nabla_{z} \phi^{j}(z', \omega), h \rangle,$   $h \in S$ 

$$\omega \in \mathbb{N}_{\delta}(\mathbb{Q}_{\varepsilon}^{\mathsf{J}}(z)), \ \mathbf{j} \in \mathbb{J}_{\varepsilon}^{\mathsf{T}}(z'); \ \mathrm{Dg}_{\varepsilon}^{\mathsf{J}}(z'), \ \mathbf{j} \in \mathbb{J}_{p}^{\mathsf{T}}\}$$

By Fact 3,  $\forall z \in \mathbb{R}^{n}$ ,  $\exists \gamma_{1} > 0$  such that  $\Omega_{\varepsilon}^{j}(z') \subset \mathbb{N}_{\delta}(\Omega_{\varepsilon}^{j}(z))$ ,  $\forall j \in J_{\varepsilon}^{f}(z')$ , So that  $\theta_{\varepsilon}(z') \leq \hat{\theta}_{\varepsilon}(z') \quad \forall z' \in \mathbb{N}_{\gamma_{1}}(z)$ .

Let  $\gamma \in (0,\gamma_1]$  be such that  $\hat{\theta}_{\epsilon}(z') \leq \frac{2}{3} \hat{\theta}_{\epsilon}(z) \neq z' \in N_{\gamma}(z)$ . Such a  $\gamma$  exists because of the continuity of  $\psi_0$ ,  $\nabla f^0$ ,  $\nabla_z \phi^j \neq j \in J_m$ ,  $\nabla g^j \neq j \in J_p$ . Hence,  $\theta_{\epsilon'}(z') \leq \hat{\theta}_{\epsilon'}(z') \leq \hat{\theta}_{\epsilon}(z') \leq \frac{2}{3} \hat{\theta}_{\epsilon}(z) \leq \frac{1}{2} \theta_{\epsilon}(z) < 0 \neq z' \in N_{\gamma}(z)$ ,  $\forall \cup \in [0,\epsilon]$ . #.

<u>Proposition 2</u>.  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\exists \gamma > 0 \exists \varepsilon > 0$  such that  $\overline{\varepsilon}(z') \ge \varepsilon \forall z' \in \mathbb{N}_{\gamma}(z)$ , where  $\overline{\varepsilon}(z)$  is the  $\varepsilon$  generated by algorithm I in Step 5 (with  $z_1 = z$ ).

<u>Proof.</u> By Fact 8,  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\exists \gamma_1 > 0$ ,  $\exists \varepsilon_1 > 0$  such that:  $\theta_{\varepsilon'}(z') \leq \frac{1}{2} \theta_{\varepsilon}(z) \stackrel{\Delta}{=} \hat{\theta} < 0 \quad \forall z' \in \mathbb{N}_{\gamma_1}(z), \quad \forall \varepsilon' \in [0, \varepsilon_1].$ 

Hence, by Fact 1,  $\overline{\theta}_{\varepsilon'}(z') \leq \hat{\theta} \quad \forall z' \in N_{\gamma_1}(z), \quad \forall \varepsilon' \in [0, \varepsilon_1].$ Let  $\delta = \min\{\varepsilon_1, -\hat{\theta}/2\}$ . By Fact 7,  $\exists \gamma \in (0, \gamma_1], \quad \exists \varepsilon \in (0, \delta]$  such that  $\tilde{\theta}_{\varepsilon'}(z') \leq \overline{\theta}_{\varepsilon'}(z') + \delta$ , i.e.  $\tilde{\theta}_{\varepsilon'}(z') \leq \hat{\theta} + \delta \quad \forall z' \in N_{\gamma}(z), \quad \forall \varepsilon' \in [0, \varepsilon].$ Now  $\hat{\theta} \leq 2\delta$  and  $\varepsilon \leq \delta \leq \varepsilon_1$  so that:

$$\overline{\theta}_{\varepsilon}(\mathbf{z}') \leq -2\delta \leq -2\varepsilon \tag{A2}$$

$$\tilde{\theta}_{c}(z') \leq -\delta \leq -\varepsilon \tag{A3}$$

 $\begin{array}{l} \forall \ z' \in \mathrm{N}_{\gamma}(z). \quad \mathrm{Since} \ \overline{\epsilon}(z') \ \mathrm{is} \ \mathrm{the} \ \mathrm{largest} \ \mathrm{number} \ \mathrm{of} \ \mathrm{the} \ \mathrm{form} \ \varepsilon_0/2^{\ell}, \ \ell \ \mathrm{a} \\ \\ \mathrm{positive} \ \mathrm{integer}, \mathrm{satisfying} \ (\mathrm{A2}) \ \mathrm{and} \ (\mathrm{A3}), \ \mathrm{the} \ \mathrm{desired} \ \mathrm{result} \ \mathrm{is} \ \mathrm{proven}. \quad \# \\ \\ \\ \underline{\mathrm{Proposition} \ 3}. \quad \forall \ z \in \mathbb{R}^n \ \mathrm{such} \ \mathrm{that} \ \theta_0(z) < 0, \ \exists \ \gamma > 0, \ \exists \ \varepsilon > 0 \ \mathrm{such} \ \mathrm{that}: \\ \\ \mathrm{f}^0(z'') \ - \ \mathrm{f}^0(z') \ \leq \ - \ \varepsilon, \quad \forall \ z' \in \mathbb{N}_{\gamma}(z) \ \cap \ \mathrm{F}, \quad \forall \ z'' \in \mathrm{A}(z'), \ \mathrm{and} \ \psi(z'') \ - \psi(z') \ \leq \ - \ \varepsilon, \\ \\ \\ \psi(z'') \ - \ \psi(z') \ \leq \ - \ \varepsilon, \quad \forall \ z' \in \mathbb{N}_{\gamma}(z) \ \cap \ \mathrm{F}^c, \quad \forall \ z'' \in \mathrm{A}(z') \ (\mathrm{where} \ \mathrm{A}: \ \mathbb{R}^n \ \rightarrow \ 2^{\mathbb{R}^n} \end{array}$ 

is defined by Steps 1 to 6 of the conceptual algorithm with "go to Step 1" in Step 6).

Proof.

By Proposition 2,  $\forall z \in \mathbb{R}^n$  such that  $\theta_0(z) < 0$ ,  $\exists \gamma_1 > 0$ ,  $\exists \varepsilon_1 > 0$ such that  $\overline{\varepsilon}(z') \geq \varepsilon_1$ , and  $\tilde{\theta}_{\overline{\varepsilon}(z')} \leq -\varepsilon_1 \quad \forall z' \in \mathbb{N}_{\gamma_1}(z)$ . (a) (upper bound on  $\psi(z'')$ ,  $z'' \in A(z')$ ). Because of the compactness of the various sets involved,  $\exists \lambda_1 \in (0,1]$ such that:

$$|\phi^{j}(z' + \lambda h, \omega) - \phi^{j}(z', \omega)| \leq \epsilon_{1}/2$$
 (A4)

 $\begin{aligned} & \forall z' \in \mathbb{N}_{\gamma_{1}}(z), \quad \forall \omega \in \Omega, \quad \forall h \in S, \quad \forall j \in J_{m}, \quad \forall \lambda \in [0, \lambda_{1}]. \\ & \text{Let } Z \stackrel{\Delta}{=} \{(j, \omega) \mid (j \notin J_{\overline{\epsilon}(z')}^{f}(z'), \quad \omega \in \Omega), \quad (j \in J_{\overline{\epsilon}(z')}^{f}(z'), \quad \omega \notin \Omega_{\overline{\epsilon}(z')}^{j}(z'))\}; \end{aligned}$ 

then  $\phi^j(z',\omega) - \psi(z') < -\varepsilon_1$ ,  $\forall (j,\omega) \in \mathbb{Z}$ . It follows from (A4) that:

$$\phi^{\mathbf{j}}(\mathbf{z'} + \lambda \overline{\mathbf{h}}(\mathbf{z'}), \omega) - \psi(\mathbf{z'}) \leq -\varepsilon_{1}/2$$
(A5)

 $\forall r' \in N_{\gamma_1}(z)$ , for all  $(j,\omega) \in Z$  where  $\overline{h}(\cdot)$  denotes  $h_{\overline{\epsilon}(\cdot)}(\cdot)$ . Also:

$$\begin{split} \phi^{\mathbf{j}}(z' + \lambda \overline{\mathbf{h}}(z'), \omega) &- \phi^{\mathbf{j}}(z', \omega) \leq \lambda [\langle \nabla_{\mathbf{z}} \phi^{\mathbf{j}}(z', \omega), \overline{\mathbf{h}}(z') \rangle \\ &+ \sqrt{n} \int_{0}^{1} \|\nabla_{\mathbf{z}} \phi^{\mathbf{j}}(z' + s\lambda \overline{\mathbf{h}}(z'), \omega) - \nabla_{\mathbf{z}} \phi^{\mathbf{j}}(z', \omega) \|_{\mathrm{ds}}] \end{split}$$

since  $\|h\| \leq \sqrt{n} \quad \forall h \in S$ . Let  $Z^{C}$  denote the complement of Z in  $J_{m} \times \Omega$ , i.e.  $Z^{C} = \{(j,\omega) \mid (j \in J_{\overline{\varepsilon}(z')}^{f}(z'), \omega \in \Omega_{\overline{\varepsilon}(z')}^{j}(z')\}$ . Now  $\langle \nabla_{z} \phi^{j}(z',\omega), \overline{h}(z') \rangle$   $\leq \tilde{\theta} \overline{\varepsilon}(z') \leq -\varepsilon_{1} \quad \forall (j,\omega) \in Z^{C}$ ; it follows from the continuity of  $\nabla_{z} \phi^{j}$ , that  $\exists \lambda_{2} \in (0,\lambda_{1}]$  such that

$$\phi^{j}(z' + \lambda \overline{h}(z'), \omega) - \psi(z') \leq -\lambda \alpha \varepsilon_{1}$$
(A6)

¥ (j,ω) ∈  $z^c$ , ¥  $z' ∈ N_{\gamma_1}(z)$ , ¥λ ∈ [0,λ<sub>2</sub>]. Combining A5, A6 we obtain:

$$f^{J}(z' + \lambda \overline{h}(z')) - \psi(z') \leq -\lambda \alpha \varepsilon_{1}$$
 (A7)

 $\forall z' \in \mathbb{N}_{\gamma_1}(z), \forall \lambda \in (0, \lambda^f], \forall j \in J_m, \text{ where } \lambda^f \in (0, \lambda_2] \text{ satisfies}$  $\lambda^f \alpha \varepsilon_1 \leq \varepsilon_1/2.$ 

Similarly, it can be shown that  $\exists \lambda^g > 0$  such that:

$$g^{\prime}(z' + \lambda \overline{h}(z')) - \psi(z') \leq -\lambda \alpha \varepsilon_{1}$$
 (A8)

 $\forall z' \in \mathbb{N}_{\gamma} (z), \ \forall \lambda \in (0, \lambda^g], \ \forall j \in J_p.$ 

Let  $\lambda_3 = \beta^k 1 > 0$ , where  $k_1$  is the smallest non-negative integer, satisfying  $\lambda_3 \leq \max{\{\lambda^f, \lambda^g\}}$ . From (A7) and (A8) and the definition of  $\psi$  we obtain:

$$\psi(z' + \lambda_{3}\overline{h}(z')) - \psi(z') \leq -\lambda_{3}\alpha\epsilon_{1}$$
(A9)

 $\forall z' \in N_{\gamma_1}(z)$ . Now, for  $z_1 = z'$ , Step 5 of the algorithm chooses the <u>smallest</u> non-negative integer  $\ell(z')$  and a corresponding  $\overline{\lambda}(z') = \beta^{\ell(z')}$  such that:

$$\psi(z'') - \psi(z') = \psi(z' + \overline{\lambda}(z')\overline{h}(z')) - \psi(z')$$

$$\leq -\overline{\lambda}(z')\alpha\varepsilon(z') \leq -\overline{\lambda}(z') \alpha \varepsilon_{1}, \qquad (A10)$$

with  $z'' = z' + \overline{\lambda}(z')\overline{h}(z')$ .

Comparing (A9) with (A10) we obtain

$$\overline{\lambda}(z') \geq \lambda_{3}$$

$$\Psi z' \in N_{\gamma_{1}}(z) \cap F^{C}. \text{ Hence:}$$

$$\psi(z'') - \psi(z') \leq -\beta\lambda_{3}\alpha\varepsilon_{1}$$

$$\Psi z' \in N_{\gamma_{1}}(z), \quad \Psi z'' \in A(z')$$
(A11)

(b) (upper bound on  $f^{0}(z'') z'' \in A(z')$ ). Let  $z' \in F(\psi_{0}(z') = 0, \psi(z') \leq 0)$ .

Since  $\psi_0(z') = 0$ , it follows that:

$$\langle \nabla f^{o}(z'), \overline{h}(z') \rangle \leq \theta_{\overline{\varepsilon}(z')}(z') \leq -\varepsilon_{1}, \forall z' \in N_{\gamma_{1}}(z).$$

Hence  $\exists \lambda_4 > 0$  such that:

$$f^{o}(z' + \lambda \overline{h}(z')) - f^{o}(z') \leq - \lambda \alpha \varepsilon_{1}$$

 $\forall \lambda \in (0,\lambda_4], \forall z' \in N_{\gamma_1}(z) \cap F.$ 

It follows from the discussion in (a) above (which is not restricted to the case  $z' \in F^{C}$ ) and, in particular, from (A9), since  $\psi(z') \leq 0$ , that  $\frac{1}{2} \lambda_{5} > 0$  such that:

$$\psi(\mathbf{z'} + \lambda \mathbf{h}(\mathbf{z'})) \leq 0$$

 $\forall z' \in \mathbb{N}_{\gamma_1}(z), \ \forall \lambda \in (0, \lambda_5]. \text{ Let } \lambda_6 = \beta^{k_2} > 0 \text{ be such that } \lambda_6 \leq \max\{\lambda_4, \lambda_5\}. \text{ If } l(z') \text{ is generated by the algorithm in Step 5, then}$  $\overline{\lambda}(z') = \beta^{l(z')} \text{ satisfies } \overline{\lambda}(z') \geq \lambda_6 \quad \forall z' \in \mathbb{N}_{\gamma_1}(z) \cap F. \text{ Hence:}$ 

$$f'(z'') - f_{o}(z') \leq -\lambda_{f} \alpha \varepsilon_{1}$$
(A13)

(A12)

#

$$\forall z' \in \mathbb{N}_{\gamma_{1}}(z), \forall z'' \in A(z').$$

(All) and (Al3) establish Proposition 3.