Quantifying the Energy Efficiency of Object Recognition and Optical Flow

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Quantifying the Energy Efficiency of Object Recognition and Optical Flow

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March 28, 2014

Abstract

In this report, we analyze the computational and performance aspects of current state-of-the-art object recognition and optical flow algorithms. First, we identify important algorithms for object recognition and optical flow, then we perform a pattern decomposition to identify key computations. We include profiles of the runtime and energy efficiency (GFLOPS/W) for our implementation of these applications on a commercial architecture. Finally, we include an analysis of memory-bandwidth boundedness for optical flow to identify opportunities for communication-avoiding algorithms.

Our results were measured on an Intel i7-4770K (Haswell) reference platform. A five-layer convolutional neural network used for object classification achieves 0.70 GFLOPS/W, which is 21% of the theoretical compute bound for this Haswell processor. On the Horn-Schunck, Lucas-Kanade, and Brox optical flow methods our implementations achieve 0.0338, 0.0103, and 0.0203 GFLOPS/W respectively. Our implementation achieves 7.9% of the theoretical bandwidth bound, assuming no cross-iteration memory optimization, for Horn-Schunck optical flow using the Jacobi solver, and 9.7% of the bandwidth bound for the conjugate-gradient solver. To improve performance, we will focus first on increasing bandwidth utilization, then on doing cross-iteration memory optimizations such as blocking and tiling the Jacobi solver and employing communication-avoiding linear solvers.

We also compare the runtime-accuracy tradeoffs for each optical flow method. We find that each method has distinct advantages over the other methods in terms of the runtime-accuracy tradeoff, so we will continue to develop and support all three methods in the future.

1 Introduction

In this report, we examine computations required for on-board unmanned aerial vehicle (UAV) vision processing. Specifically, we focus on object recognition, object tracking, and optical flow. Given that on-board processing is constrained by power, we focus on quantifying the energy efficiency and accuracy of current state-of-the-art methods.

We start by decomposing the application capabilities, for example object recognition using convolutional neural networks, into patterns. This provides a high-level structural and computational understanding of the application. Then we profile performance on a commercial processor (Intel i7-4770K). To contextualize performance, we calculate the number of floating point operations (FLOPs) performed for each computation. This allows us to compute a measure of energy efficiency known as giga-flops per second per watt (GFLOPS/W). For optical flow, we also count the total number of bytes transferred in the inner loop of the algorithm. This
Figure 1: Deep convolutional neural network for object recognition [2]. This network layout achieved the highest object classification accuracy for ImageNet 2012, and a similar network won ImageNet 2013. Big data and efficient implementations have made deep learning accurate and tractable.

allows us to compute the flop-to-byte ratio, which gives an idea of the memory-boundedness of the algorithm and the potential speedup from communication-avoiding algorithms [1].

In Section 2, we analyze object recognition using convolutional neural networks [2]. We find that a five-layer convolutional neural network is able to achieve 0.70 GFLOPS/W. In Section 3, we analyze the Horn-Schunck, Lucas-Kande, and Brox optical flow algorithms [3, 4, 5]. We determine these methods achieve 0.0338, 0.0103, and 0.0203 GFLOPS/W respectively. We also compute the flop-per-byte ratio for Horn-Schunck and Brox methods and report accuracy results for all three methods, and we compute the achieved percentage of peak bandwidth for the Horn-Schunck method. Section 4 concludes our report.

2 Object Recognition

Object recognition is a key enabling technology for a variety of UAV capabilities including navigation, odometry, and reconnaissance.

2.1 State of the art Algorithm

Within the past 18 months, the computer vision community has seen a large improvement in accuracy by designing systems based on deep neural networks instead of hand-engineered descriptors. The key algorithms of the deep learning revolution can be traced back to the late 1980s. However, the rise of big data has led to huge labeled datasets (e.g. ImageNet [6] with >1M labeled images) for training and evaluating object recognition systems. It turns out that large datasets are a lynchpin of high-accuracy neural networks for object recognition. Additionally, extremely efficient deep neural network implementations such as Berkeley’s Caffe [7] expose enough parallelism to make ImageNet a tractable benchmark for deep neural network object recognition. Today, neural networks such as Alexnet [2] and their ilk (e.g. [8],[9]) provide state-of-the-art object classification accuracy (up to 88% when scored on top-5 categories) on the 1000-category ImageNet dataset. We show an illustration of Alexnet in Figure 1.

In the remainder of this section, we analyze the computational patterns and bottlenecks, GFLOPS/s, and energy for state-of-the-art deep convolutional networks on the Haswell reference architecture.
2.2 Computational Patterns

After several years of work on efficient computer vision in the ParLab and ASPIRE Lab at Berkeley, we have codified computer vision computations into a “periodic table” of 15 underlying computational patterns (Figure 2). While computer vision algorithms continue to evolve and advance, these underlying patterns have remained relatively static over many generations of computer vision algorithms. Recently, object recognition algorithms have seen a major shift to deep learning, and it would be easy for computational efficiency researchers to be intimidated by this – how much of what we know about efficient computer vision will transfer to these new deep neural algorithms? Well, as it turns out, it is quite reasonable to map deep neural networks into our periodic table of computer vision patterns (Figure 2). As we will show in the next paragraph, analyzing the performance and energy efficiency of deep neural networks is quite easy, so long as we think in terms of well-understood patterns that are less susceptible to computer vision algorithmic changes.

Broadly, deep neural networks perform feature extraction and recognition by taking an image and feeding it through several layers of filters and dimensionality adjustments. For both training and inference, layers are implemented with primitives such as 3D convolution with multiple kernels, neighborhood max filtering (“max-pooling”), ReLu (removing negative numbers), and dropout (zeroing out a random collection of values to avoid overfitting). In Table 1, we map these deep neural network primitives into our periodic table of computer vision computational patterns.

2.3 Performance and Energy Analysis

We now turn to analyzing the computational complexity, efficiency, and energy of object recognition with convolutional neural networks. As shown in the previous subsection, convolution dominates the overall computation time in this system. Therefore, we use convolution as a lower bound in terms of overall FLOP count, leading to slightly conservative but reasonable
Table 1: Mapping convolutional neural networks to the ASPIRE periodic table of computer vision patterns.

<table>
<thead>
<tr>
<th>Layer type</th>
<th>Pattern(s)</th>
</tr>
</thead>
<tbody>
<tr>
<td>convolution layer</td>
<td>convolution</td>
</tr>
<tr>
<td>dropout layer</td>
<td>–</td>
</tr>
<tr>
<td>ReLu layer</td>
<td>convolution (1x1 filter)</td>
</tr>
<tr>
<td>max-filtering</td>
<td>convolution-style data access pattern</td>
</tr>
</tbody>
</table>


<table>
<thead>
<tr>
<th>Layer</th>
<th>Runtime per 50 frames (s)</th>
<th>Input dims</th>
<th>Filter dims</th>
<th># filters</th>
<th>Complexity (# GFLOP)</th>
<th>GFLOPS/s</th>
<th>Avg Power (W)</th>
<th>Energy per 50 frames (J)</th>
<th>GFLOPS/s/W</th>
</tr>
</thead>
<tbody>
<tr>
<td>conv1</td>
<td>0.0980</td>
<td>224x224x3</td>
<td>11x11x3</td>
<td>96</td>
<td>10.93</td>
<td>111.46</td>
<td>66.97</td>
<td>&gt;=127.28</td>
<td>0.70</td>
</tr>
<tr>
<td>conv2</td>
<td>0.1806</td>
<td>55x55x48</td>
<td>5x5x48</td>
<td>256</td>
<td>22.39</td>
<td>124.02</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>conv3</td>
<td>0.0860</td>
<td>27x27x128</td>
<td>3x3x128</td>
<td>384</td>
<td>14.95</td>
<td>173.81</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>conv4</td>
<td>0.0850</td>
<td>13x13x192</td>
<td>3x3x192</td>
<td>384</td>
<td>11.21</td>
<td>130.52</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>conv5</td>
<td>0.0756</td>
<td>13x13x192</td>
<td>3x3x192</td>
<td>256</td>
<td>7.48</td>
<td>98.93</td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>TOTAL conv layers</td>
<td>0.5261</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
<tr>
<td>TOTAL all layers</td>
<td>0.71</td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

In Table 2, we show the analysis and results in terms of GFLOPS/s for the convolutional layers in the Caffe convolutional neural network, using the same configuration as discussed previously. As you can see in Table 2, the layers funnel down from a 224x224 3-channel input image down to a 13x13 256-channel feature descriptor map. Given our coarse power measurement technology, we are able to obtain the power and energy of the overall system, but not of individual layers. Also, since the computation of neural networks in Caffe is dominated by convolution, we use the number of GFLOPs in convolution as a lower bound for the overall computational complexity of all layers in the neural network. This analysis culminates in finding that object recognition with the Caffe convolutional neural network achieves 0.70 GFLOPS/s/W on the Intel Haswell reference platform (Table 2). This is 21% of the theoretical compute bound for this Haswell processor.

3 Optical Flow

Optical flow is a common computer vision application that computes the apparent motion of each pixel between pairs of images, or between frames in a video. Optical flow information enables point tracking which can be a powerful capability for UAVs.

Optical flow between two images can be visualized (Figure 3b). Colors indicate direction and intensity indicates magnitude of pixel motion. Quality of solution is measured using standard benchmarks such as the Middlebury optical flow benchmark dataset [10] and the KTTI vision benchmark suite [11]. The quality metrics for optical flow are average angular error (AAE) of the flow vectors compared to ground truth provided by the these benchmark datasets, as well as average endpoint error (AEE) of the flow vectors.

There are many different ways to solve optical flow. As of February 2014, the KTTI vision benchmark results webpage for optical flow reports results for 42 different optical flow methods. The Middlebury benchmark results webpage reports results for 95 different optical flow methods.

We choose three methods to focus on: Horn-Schunck [3], Brox [5], and Lucas-Kanade [12]. We focus on Horn-Schunck and Brox due to their popularity, along with the general consensus.
that the majority of newer methods are simply extensions of these original formulations [13].
We include Lucas-Kanade in our analysis because it considers only local image patches, so it is fundamentally different than Horn-Schunck and Brox. The Lucas-Kanade method is also much cheaper to compute.

We will analyze these three optical flow methods in the following subsections. For each method, as specified in the deliverable text, we will show:

- the high level algorithm description,
- the decomposition into computational and structural patterns,
- profiles of the runtime and energy analysis,
- analysis of computation vs. quality of solution trade-offs, and
- an analysis of memory-boundedness.

We also provide plots comparing all three optical flow methods side-by-side, run with a variety of parameters, in terms of runtime vs. accuracy.

3.1 Horn-Schunck Method

3.1.1 High level algorithm description

The Horn-Schunck method is formulated as a minimization of the following energy functional:

\[
E = \iint (I_x u + I_y v + I_t)^2 + \alpha^2 (\|\nabla u\|^2 + \|\nabla v\|^2) dxdy
\]  

(1)

\(I_x\) is the image gradient in the \(x\) dimension, \(I_y\) is the image gradient in the \(y\) dimension, \(I_t\) is the image gradient in the time dimension, \(u\) and \(v\) are the \(x\) and \(y\) components of the flow
def hs_oflow (im1_data, im2_data, # Input images
    D, # Laplacian stencil
    Gx, Gy, # Gradient stencils
    u, v, # Flow vectors
    zero, one, lam2 # Scalars
):
    du = zero * u # Dense Linear Algebra
    dv = zero * v # Dense Linear Algebra
    Ix = Gx*im1_data # Structured Grid
    Iy = Gy*im1_data # Structured Grid
    It = im1_data -
        warp_img2d(im2_data, u, v) # Sparse Linear Algebra
    Ix2 = Ix * Ix # Dense Linear Algebra
    IxIy = Ix * Iy # Dense Linear Algebra
    Iy2 = Iy * Iy # Dense Linear Algebra
    # Application pattern: Linear Solver
    # Structural pattern: Iterator
    for i in range(200):
        ubar = D * du # Structured Grid
        vbar = D * dv # Structured Grid
        num = Ix * ubar + Iy * vbar + It # Dense Linear Algebra
        den = Ix2 + Iy2 + lam2 # Dense Linear Algebra
        du = ubar - (Ix * num) / den # Dense Linear Algebra
        dv = vbar - (Iy * num) / den # Dense Linear Algebra
    return du, dv

Figure 4: Code and pattern decomposition for Horn-Schunck optical flow solved using the Jacobi linear solver

vectors, respectively, and $\alpha$ is a parameter that trades off smoothness of the flow field with the accuracy of the flow nearby a given pixel. The integral is summing this quantity for every pixel in the image. Since each term in this functional represents a cost, we try to minimize the functional over all possible flow functions $u$ and $v$. A detailed derivation of the Horn-Schunck algorithm is included in the appendices of this document.

3.1.2 Decomposition into computational and structural patterns

Figure 4 contains code that computes optical flow in Python using the Horn-Schunck method and the Jacobi linear solver. This code runs in our Hindemith framework. The Hindemith framework analyzes the algorithm description in Python and composes hand-written OpenCL functions to eliminate unnecessary memory traffic that is common array codes such as this. This particular implementation was designed to match the implementation found in the ArrayFire example codes [14]. We’ve annotated the code with the input types, as well as the computational and structural patterns.
<table>
<thead>
<tr>
<th>Solver Type</th>
<th>Runtime per Frame (s)</th>
<th>Average Power (W)</th>
<th>Energy per Frame (J)</th>
<th># Iter.</th>
<th>Runtime per iteration (ms)</th>
<th>Energy per iteration (mJ)</th>
<th>GFLOPS per Watt</th>
</tr>
</thead>
<tbody>
<tr>
<td>Jacobi</td>
<td>0.735</td>
<td>95.16</td>
<td>69.98</td>
<td>400</td>
<td>1.838</td>
<td>174.9</td>
<td>0.0338</td>
</tr>
<tr>
<td>CG</td>
<td>1.542</td>
<td>98.39</td>
<td>151.74</td>
<td>400</td>
<td>3.856</td>
<td>379.4</td>
<td>0.0256</td>
</tr>
<tr>
<td>PCG</td>
<td>2.193</td>
<td>100.69</td>
<td>220.80</td>
<td>400</td>
<td>5.482</td>
<td>552.0</td>
<td>0.0252</td>
</tr>
<tr>
<td>RB</td>
<td>2.539</td>
<td>96.43</td>
<td>244.79</td>
<td>400</td>
<td>6.346</td>
<td>612.0</td>
<td>0.0193</td>
</tr>
</tbody>
</table>

Table 3: Runtime and energy metrics for Horn-Schunck optical flow

The solver consists mainly of Dense Linear Algebra and Structured Grid operations. This pattern decomposition informs hardware and software implementation choices. For example, these particular patterns are particularly amenable to vectorization and tiling optimizations so we expect this application to compile to efficient code using vectorizing compilers or implicitly parallel languages, and to execute efficiently on vector or SIMD hardware.

We can also solve the Horn-Schunck algorithm using other linear solvers, such as conjugate-gradient (CG), preconditioned conjugate-gradient (PCG), and red-black Gauss Seidel (RB). We have also implemented those linear solvers in our Hindemith framework and will present comparative performance and accuracy results using those linear solvers.

3.1.3 Profiles of runtime and energy analysis

In this section we will analyze the runtime and energy consumption of Horn-Schunck optical flow using a variety of different linear solvers. We run our experiment on Intel Core i7-4770 CPU 3.4 GHz (Haswell) processor. We use the AMD APP OpenCL SDK compiler and runtime. The power was taken using the Watts Up Pro? power logger at one second intervals, then averaged over the duration of multiple executions. The input image pair is RubberWhale from the Middlebury optical flow benchmark set, which is size 588x384. We do not resize the image. We represent the images, the flow, and all intermediate data in grayscale single-precision floating point format. We set the parameter $\alpha$ to 0.1.

Because the Hindemith framework fuses most of these operations together, we do not have line-by-line profiling information. Instead, we record the runtime and power consumption for an execution of 400 linear solver iterations. Then we compute the average runtime and energy per linear solver iteration. Some linear solvers converge faster than others. We will consider the trade-off between per-iteration efficiency and convergence rate in the next section.

3.1.4 Analysis of computation vs quality of solution trade-offs

As shown in Table 3, different linear solvers have different runtime and energy costs. However, there is a trade-off between computation and quality of solution that must be explored. More computationally expensive solvers such as conjugate-gradient converge to a good solution faster than cheaper solvers like Jacobi. Figure 5 shows the rate of convergence for each of the four linear solvers. The best performing solver, in terms of quality per number of iterations, is the preconditioned conjugate gradient solver, followed by the red-black Gauss Seidel, and conjugate gradient solvers.

We are interested in the energy it takes to get a given solution quality. In Figure 6, we plot the quality of solution for each linear solver per Joule. This is calculated by multiplying
the average number of Joules per iteration with the total number of iterations at each point. As we would expect, Jacobi becomes a more attractive option when we consider Joules instead of iterations. This is because Jacobi iterations are comparatively cheap in terms of Joules. Conversely, our implementation of red-black Gauss Seidel proves to be very inefficient when we consider Joules instead of just iterations.

3.1.5 Analysis of memory-boundedness

For our analysis of memory-boundedness and GFLOPS/Watt, we consider only the inner loop of the linear solver. In Figure 7, we annotate the number of FLOPS performed for every operation in the Jacobi linear solver:

One iteration of the Jacobi linear solver for Horn-Schunck computes $28h^2w$ FLOPS, where $h$ and $w$ are the height and width of the image respectively. The number of words transferred between memory and the processor is $9h^2w$ (to read variables $I_x$, $I_y$, $I_t$, $I_{x2}$, $I_{y2}$, $d_u$, $d_v$, and to write variables $d_u$, $d_v$). We can use this information to compute the arithmetic intensity (FLOPS/Byte) of this kernel. Arithmetic intensity is a measure of the memory-boundedness of a particular algorithm or implementation [15]. Depending on the balance of floating-point and bandwidth capabilities of a particular device, we can compute a limit on the achievable performance by multiplying the arithmetic intensity by the peak memory bandwidth, computed by the STREAM benchmark [16].

The arithmetic intensity of the Jacobi optical flow kernel is $\frac{28}{429} = 0.778$ FLOPS per byte. We measured a STREAM copy bandwidth of up to 52 GB/Sec on the Haswell reference machine with an array size set to match the working set of the Jacobi solver (compiler flags: gcc -O3 -fopenmp stream.c -o stream_omp -DSTREAM_ARRAY_SIZE=396536, 8 threads). This means the max achievable performance for this algorithm on the Haswell reference platform, without cross-iteration memory optimizations, is $0.778 \times 52$, or 40.4 GFLOPS. Our implementation actually performs $13h^2w$ word transfers instead of the $9h^2w$ theoretical minimum because the framework breaks the inner loop into two OpenCL kernels. We are currently achieving 3.21 GFLOPS, which is 8% of the 40.4 GFLOPS limit imposed by the bandwidth-bound, assuming
Figure 6: Quality of solution vs number of Joules spent for different Horn-Schunck linear solvers

Figure 7: Number of FLOPS performed for each operation in the Horn-Schunck Jacobi linear solver

ubar = D * du  # 9*h*w FLOPS
vbar = D * dv  # 9*h*w FLOPS
num = Ix * ubar + Iy * vbar + It  # 4*h*w FLOPS
den = Ix2 + Iy2 + lam2  # 2*h*w FLOPS
du = ubar - (Ix * num) / den  # 2*h*w FLOPS
dv = vbar - (Iy * num) / den  # 2*h*w FLOPS
Ap0 = (D*p0 + Ix2 * p0 + IxIy * p1) # 13*h*w FLOPS
Ap1 = (D*p1 + Iy2 * p1 + IxIy * p0) # 13*h*w FLOPS
alpha = rsold / sum2d(p0 * Ap0 + p1 * Ap1) # 4*h*w FLOPS
du = du + alpha * p0 # 2*h*w FLOPS
dv = dv + alpha * p1 # 2*h*w FLOPS
r0 = r0 - alpha * Ap0 # 2*h*w FLOPS
r1 = r1 - alpha * Ap1 # 2*h*w FLOPS
rsnew = sum2d(r0 * r0 + r1 * r1) # 4*h*w FLOPS
beta = rsnew / rsold
p0 = r0 + beta * p0 # 2*h*w FLOPS
p1 = r1 + beta * p1 # 2*h*2 FLOPS
rsold = rsnew

Figure 8: Number of FLOPS performed for each operation in the Horn-Schunck conjugate gradient linear solver

no cross-iteration memory optimizations. In the future, we can work to improve our bandwidth performance. The application can also benefit from the cross-iteration memory optimizations. This includes blocking and tiling the Jacobi solver and employing communication-avoiding linear solvers.

Figure 8 shows the number of FLOPS computed for the Horn-Schunck method using the conjugate-gradient solver. One iteration of the conjugate-gradient linear solver for Horn-Schunck computes 46*h*w FLOPS, where h and w are the height and width of the image respectively. The number of words transferred between memory and the processor is 15*h*w (to read variables du, dv, p0, p1, Ix2, IxIy, Iy2, r0, r1, and to write variables p0, p1, r0, r1, du, dv). This means the arithmetic intensity of this kernel is \(\frac{46}{4*15} = 0.766\) FLOPS per byte. We measured a STREAM copy bandwidth of up to 33 GB/Sec on the Haswell reference machine with an array size set to match the working set of the Jacobi solver (compiler flags: gcc -O3 -fopenmp stream.c -o stream -DSTREAM_ARRAY_SIZE=509832, 8 threads). This means the limit imposed by the bandwidth-bound for this algorithm is 0.766 \times 33 = 25.7 GFLOPS. We are currently achieving 2.51 GFLOPS, or 9.7% of the limit. In the future, we can work to improve our bandwidth performance. After optimizing for bandwidth performance, we can investigate using communication-avoiding linear solvers, which can reduce both the amount of data transferred and the number of synchronizations performed.

We compute the number of GFLOPS for the rest of the Horn-Schunck linear solvers using a similar approach. This allows us to compute the GFLOPS/W values in Table 3.

3.2 Lucas-Kanade Method

3.2.1 Problem Formulation

The Lucas-Kanade method is works by examining local regions across a pair of image and computing a displacement vector for each local region. The method can be visualized using a figure from the original 1981 paper [12], shown here as Figure 9.

The displacement vector for each local region is computed using a least squares solution to the optical flow equation (derived in Appendix A) solved simultaneously for each pixel in the region:
Figure 9: From the original Lucas-Kanade paper (1981): “We wish to find the disparity vector \( h \) which minimizes some measure of the difference between \( F(x + h) \) and \( G(x) \) for \( x \) in some region of interest \( R \)".

![Diagram of Lucas-Kanade algorithm](image)

<table>
<thead>
<tr>
<th>Runtime per Frame (s)</th>
<th>Average Power (W)</th>
<th>Energy per Frame (J)</th>
<th># Iter.</th>
<th>Runtime per iteration (ms)</th>
<th>Energy per iteration (mJ)</th>
<th>GFLOPS per Watt</th>
</tr>
</thead>
<tbody>
<tr>
<td>0.02286</td>
<td>59.84</td>
<td>1.368</td>
<td>1</td>
<td>22.86</td>
<td>1368</td>
<td>0.0103</td>
</tr>
</tbody>
</table>

Table 4: Runtime and energy metrics for Lucas-Kanade optical flow using a 4x4 window size

\[
\frac{\partial I}{\partial x} u + \frac{\partial I}{\partial y} v + \frac{\partial I}{\partial t} = 0 \tag{2}
\]

The Lucas-Kanade algorithm iterates solving this least squares problem and warping the image until a local minimum is found. This process is described succinctly in a review of Horn-Schunck by Baker and Matthews [4].

### 3.2.2 Pattern Decomposition

Figure 10 shows our implementation of Horn-Schunck that runs in Python, as well as our Hindemith framework. We’ve annotated the code with the input types, as well as the computational and structural patterns within the application.

Like Horn-Schunck, this method consists of mainly the Dense Linear Algebra and Structured Grid computational patterns. This tells us that the application will be amenable to optimizations such as tiling and vectorization. However, there is a data-dependent array indexing operation, image warping, that falls under the Sparse Linear Algebra pattern. This means that we won’t be able to statically partition the computation given that dependences between operations are not known until runtime.

### 3.2.3 Profiles of runtime and energy analysis

We evaluate the runtime and energy performance of Lucas-Kanade optical flow using the same machine, methodology, and input data as was used for the Horn-Schunck analysis in Section 3.1.3. In this case, we only run one iteration of Lucas-Kanade and our window size is 4x4. The energy per iteration is much higher than Horn-Schunck, but the energy per problem is much lower.
def lk_oflow (I1, I2, # Input images
           Ix, Iy, # Structured Grid
           Gx, Gy): # Gradient stencils
    Ix = Gx * I2 # Structured Grid
    Iy = Gy * I2 # Structured Grid

    # Structural pattern: Iterator
    for i in range(2):
        WarpedI2 = warp_img2d(I2, u, v) # Sparse Linear Algebra
        WarpedIx = warp_img2d(Ix, u, v) # Sparse Linear Algebra
        WarpedIy = warp_img2d(Iy, u, v) # Sparse Linear Algebra
        ErrorImg = I1 - WarpedI2 # Dense Linear Algebra

        du, dv = lk_least_squares(ErrorImg, WarpedIx, WarpedIy)

        u = u + du # Dense Linear Algebra
        v = v + dv # Dense Linear Algebra

    return u, v

Figure 10: Code and pattern decomposition for Lucas-Kanade optical flow

3.2.4 Analysis of computation vs quality of solution trade-offs

Figure 11 shows the quality of solution (average angular error) of the Lucas-Kanade method for our benchmark image. The average angular error for the Lucas-Kanade method is lower than both the Horn-Schunck and the Brox methods for the same image. However, the Lucas-Kanade solution is much cheaper to compute.

3.2.5 Analysis of memory-boundedness

The Lucas-Kanade algorithm is dominated by the number of FLOPs required for the solution to the local least squares problems. This depends on the size of the window under consideration. For larger windows, the computation will likely be compute-bound, for very small windows the computation will most likely be memory-bound.

We run our tests with a 4x4 window. In this case, the number of FLOPS computed per pixel is 2*2*winsize*winsize, plus 3 FLOPS for computing the error image and adding the displacements to the flow vectors. So the total number of FLOPS for our example is 67*h*w, where h and w are the height and width of the image. This FLOP count is used to compute GFLOPS per Watt in Table 4.

3.3 Brox Method

3.3.1 Problem Formulation

The Brox method is a recent algorithm for optical flow that attains high-quality results but also comparitively computationally intensive [5]. It is based on an energy-minimization approach.
Figure 11: Quality of solution for multiple Lucas Kanade iterations using a 4x4 window size similar to the Horn-Schunck method. However, the functional to be minimized has additional terms that add complexity to the formulation but produce a higher-quality result.

The main optimization problem is to minimize an energy functional with two terms. The first term enforces adherence to the gray-value constancy and gradient constancy assumptions, and the second term enforces smoothness in the flow field. In these equations we use the notation from the original paper, where the image coordinates (both x and y) are represented by a vector \( x \), and the flow field (both u and v) are represented by a vector \( u \):

\[
E(u, v) = E_{\text{Data}} + \alpha E_{\text{Smooth}}
\]

The \( E_{\text{Data}} \) term penalizes deviations from the gray-value constancy assumption (described in Appendix A), and the gradient constancy assumption:

\[
E_{\text{Data}}(u, v) = \int_{\Omega} \left( |I(x + w) - I(x)|^2 + \gamma |\nabla I(x + w) - \nabla I(x)|^2 \right) dx
\]

The \( E_{\text{Smooth}} \) term penalizes flow functions that are not smooth. That is, typically pixels that are nearby one another move in similar directions and at similar speeds. So we assume the change in flow should not vary much in either space or time:

\[
E_{\text{Smooth}}(u, v) = \int \Psi(\|\nabla_3 u\|^2 + \|\nabla_3 v\|^2) dx
\]

The \( \Psi \) function is added to this penalty function in order to reduce the influence of outliers. \( \epsilon \) is 1E-3:

\[
\Psi(s^2) = \sqrt{s^2 + \epsilon^2}
\]

3.3.2 Pattern Decomposition

The following code computes optical flow in Python using the Brox method and the conjugate-gradient linear solver. This particular implementation was designed to match the implementation found in the OpenCV library [17]. However, this code does run in our Hindemith framework. We’ve annotated the code with the input types, as well as the computational and structural patterns.

```python
def brox_oflow (I1, I2, # Input images
               Gx, Gy, # Gradient stencils
```

u, v, pointfive, zero, eps, brox_alpha, brox_beta, gamma, A

\begin{verbatim}
du = zero * u
dv = zero * v
tex_Ix0 = Gx * I1

tex_Iy0 = Gy * I1

tex_Ix = Gx * I2

tex_Iy = Gy * I2

tex_Ixx = Gx * tex_Ix

tex_Iyy = Gy * tex_Iy

tex_Ixy = Gx * tex_Iy

# Application pattern: Non-convex non-linear solver
# Outer fixed-point iterations
for outer in range(5):

    # Warp images
    Iz = warp_img2d(I2, u, v) - I1
    Ix = warp_img2d(tex_Ix, u, v)
    Ixz = Ix - tex_Ix0
    Ixy = warp_img2d(tex_Ixy, u, v)
    Ixx = warp_img2d(tex_Ixx, u, v)
    Iy = warp_img2d(tex_Iy, u, v)
    Iyz = Iy - tex_Iy0
    Iyy = warp_img2d(tex_Iyy, u, v)

    pd1 = (Iz + Ix * du + Iy * dv)
    pd2 = (Ixz + Ixx * du + Ixy * dv)
    pd3 = (Iyz + Ixy * du + Iyy * dv)

    PsiData = pointfive / 
    sqrt(pd1*pd1 + gamma*(pd2*pd2 + pd3*pd3) + eps)

    # Set up linear system
    gx0 = Gx * (u + du)
    gy0 = Gy * (u + du)
    gx1 = Gx * (v + dv)
    gy1 = Gy * (v + dv)
    PsiSmooth = pointfive / 
    sqrt(gx0*gx0 + gy0*gy0 + gx1*gx1 + gy1*gy1 + eps)

    set_brox_matrix(A, PsiSmooth, brox_alpha) 

    du_coef0 = PsiData * (Ix * Ix + gamma * (Ixx * Ixx + Ixy * Ixy))
    dv_coef0 = PsiData * (Ix * Iy + gamma * (Ixx * Ixy + Ixy * Iyy))
    du_coef1 = PsiData * (Iy * Ix + gamma * (Iyy * Ixy + Ixy * Ixx))
    dv_coef1 = PsiData * (Iy * Iy + gamma * (Iyy * Iyy + Ixy * Ixy))
    b0 = zero - (PsiData * Ix + gamma * PsiData * 
               (Ixx * Ix + Ixy * Iy) + A * (u + du))
    b1 = zero - (PsiData * Iy + gamma * PsiData * 
               (Iyy * Ixy + Iy + Ixz) + A * (v + dv))
\end{verbatim}
<table>
<thead>
<tr>
<th>Solver Type</th>
<th>Runtime per Frame (s)</th>
<th>Average Power per Frame (W)</th>
<th>Energy per Frame (J)</th>
<th># Iter.</th>
<th>Runtime per Iteration (ms)</th>
<th>Energy per Iteration (mJ)</th>
<th>GFLOPS per Watt</th>
</tr>
</thead>
<tbody>
<tr>
<td>CG</td>
<td>6.527</td>
<td>101.4</td>
<td>662.1</td>
<td>1400</td>
<td>4.662</td>
<td>472.9</td>
<td>0.0202</td>
</tr>
<tr>
<td>PCG</td>
<td>8.656</td>
<td>104.2</td>
<td>901.8</td>
<td>1400</td>
<td>6.183</td>
<td>644.1</td>
<td>0.0203</td>
</tr>
</tbody>
</table>

Table 5: Runtime and energy metrics for Brox optical flow

```python
# Application pattern: Linear system of equations
# Solve CG for du, dv
r0 = b0 - (A*du + du_coef0 * du + dv_coef0 * dv) # Structured Grid
r1 = b1 - (A*dv + du_coef1 * du + dv_coef1 * dv) # Structured Grid
p0 = b0 - (A*du + du_coef0 * du + dv_coef0 * dv) # Structured Grid
p1 = b1 - (A*dv + du_coef1 * du + dv_coef1 * dv) # Structured Grid
rsold = sum2d(r0 * r0 + r1 * r1) # Dense Linear Algebra
for inner in range(40):
    Ap0 = (A * p0 + du_coef0 * p0 + dv_coef0 * p1) # Structured Grid
    Ap1 = (A * p1 + du_coef1 * p0 + dv_coef1 * p1) # Structured Grid
    alpha = rsold / sum2d(p0 * Ap0 + p1 * Ap1) # Dense Linear Algebra
    du = du + alpha * p0 # Dense Linear Algebra
    dv = dv + alpha * p1 # Dense Linear Algebra
    r0 = r0 - alpha * Ap0 # Dense Linear Algebra
    r1 = r1 - alpha * Ap1 # Dense Linear Algebra
    rsnew = sum2d(r0 * r0 + r1 * r1) # Dense Linear Algebra
    beta = rsnew / rsold # Dense Linear Algebra
    p0 = r0 + beta * p0 # Dense Linear Algebra
    p1 = r1 + beta * p1 # Dense Linear Algebra
    rsold = rsnew
return du, dv
```

Like Horn-Schunck, Brox is mainly Dense Linear Algebra and Structured Grid operations which means it will run well on architectures that support vector or SIMD execution.

We can also solve the Brox algorithm using other linear solvers, such as preconditioned conjugate-gradient (PCG) and red-black Gauss Seidel. We have also implemented preconditioned conjugate-gradient in our Hindemith framework and will present comparative performance and accuracy results using both linear solvers.

3.3.3 Profiles of runtime and energy analysis

We evaluate the runtime and energy performance of Brox optical flow using the same machine, methodology, and input data as was used for the Horn-Schunck analysis in Section 3.1.3. In this case, we do seven warping iterations on the image, five outer iterations and 40 linear solver iterations, for a total of 1400 linear solver iterations. We set the parameter $\alpha$ to 0.197 and the parameter $\gamma$ to 50.0. The energy per iteration is for Brox is slightly more than the energy per iteration for Horn-Schunck even though it is the same linear solver. This is because Brox uses an explicit sparse matrix, requiring more memory traffic.
\[
\begin{align*}
\text{Ap}_0 &= (A \times p_0 + du\_coef0 \times p_0 + dv\_coef0 \times p_1) \quad \# 13h\times w \text{ FLOPS} \\
\text{Ap}_1 &= (A \times p_1 + du\_coef1 \times p_0 + dv\_coef1 \times p_1) \quad \# 13h\times w \text{ FLOPS} \\
\alpha &= \frac{rsold}{\text{sum2d}(p_0 \times \text{Ap}_0 + p_1 \times \text{Ap}_1)} \quad \# 4h\times w \text{ FLOPS} \\
\text{du} &= \text{du} + \alpha \times p_0 \quad \# 2h\times w \text{ FLOPS} \\
\text{dv} &= \text{dv} + \alpha \times p_1 \quad \# 2h\times w \text{ FLOPS} \\
\text{r}_0 &= \text{r}_0 - \alpha \times \text{Ap}_0 \quad \# 2h\times w \text{ FLOPS} \\
\text{r}_1 &= \text{r}_1 - \alpha \times \text{Ap}_1 \quad \# 2h\times w \text{ FLOPS} \\
\text{rsnew} &= \text{sum2d}(r_0 \times r_0 + r_1 \times r_1) \quad \# 4h\times w \text{ FLOPS} \\
\beta &= \frac{\text{rsnew}}{rsold} \\
\text{p}_0 &= \text{r}_0 + \beta \times p_0 \quad \# 2h\times w \text{ FLOPS} \\
\text{p}_1 &= \text{r}_1 + \beta \times p_1 \quad \# 2h\times w \text{ FLOPS} \\
\text{rsold} &= \text{rsnew}
\end{align*}
\]

Figure 12: Number of FLOPS computed for each operation in the Brox conjugate-gradient linear solver

### 3.3.4 Analysis of computation vs quality of solution trade-offs

Currently our Brox implementation achieves an average angular error of 7.33 when solved with the preconditioned conjugate-gradient algorithm using seven warping iterations, five fixed-point iterations, and 40 PCG iterations (1400 total iterations). For traditional CG with the same setup, we are currently achieving an average angular error of 8.89. As shown in Figures 13 and 14, these are among the highest accuracy results we were able to produce on our two benchmarks. It is difficult to contextualize the quality of these solutions compared to published figures because our two images come from the Middlebury training set and most published results report accuracy on the test set. The Middlebury optical flow benchmark webpage lists accuracy results for many published methods on the test set.

### 3.3.5 Analysis of memory-boundedness

The following analysis is for the Brox method using the conjugate-gradient solver. We consider the inner loop of the conjugate-gradient solver, shown in Figure 12. The only difference between this and the Horn-Schunck conjugate-gradient analysis is that the matrix must be represented explicitly, meaning it requires 5h\times w words to be transferred from memory to the processor. The Horn-Schunck sparse matrix, in contrast, could be represented implicitly.

One iteration of the conjugate-gradient linear solver for Brox computes 46h\times w FLOPS, where h and w are the height and width of the image respectively. The number of words transferred between memory and the processor is 20h\times w (to read variables A, du, dv, p0, p1, Ix2, IxIy, Iy2, r0, r1, and to write variables p0, p1, r0, r1, du, dv). This means the arithmetic intensity of this kernel is \( \frac{46}{20} = 0.575 \) FLOPS per byte. Since this solver is similar to Horn-Schunck we will take a similar approach to optimizing it: First maximize bandwidth utilization, then apply cross-iteration optimizations and communication-avoiding algorithms [1].

### 3.4 Comparison of Optical Flow Methods

In Figures 13 and 14, we compare the performance of all three optical flow methods and the linear solvers in terms of both accuracy and runtime on two Middlebury benchmark images [10].
Figure 13: Comparison of optical flow methods. Accuracy vs. runtime on the Dimetrodon benchmark image using many different configurations. Run on AMD Radeon 7990.

Figure 14: Comparison of optical flow methods. Accuracy vs. runtime on the RubberWhale benchmark image using many different configurations. Run on AMD Radeon 7990.
Figure 13 shows the accuracy/speed tradeoff for the Dinetron benchmark image, and Figure 14 shows the accuracy/speed tradeoff for the RubberWhale image. These results are on an AMD Radeon 7990 GPU, our fastest platform. We run a number of different parameterizations of each optical flow method. We vary the number of solver iterations, number of warping iterations, number of pyramid levels, and linear solver type. Each point in the plot represents a single run. For Lucas-Kanade, we vary the radius between 3, 5, and 7 pixels wide. For Horn-Schunck we set $\alpha$ to 0.1. For the Brox method, we set $\alpha$ to 0.197 and $\gamma$ to 50.0. Finally, we apply a $5 \times 5$ median filter between warping iterations and pyramid levels to improve accuracy.

The fastest, least accurate points belong to the Lucas Kanade method. As we would expect, the Brox method is generally the most accurate and the slowest. The Horn-Schunck method lies somewhere in between. The best performing linear solver varies depending on both the benchmark image and the error metric.

4 Conclusion

We’ve explored state-of-the-art algorithms and approaches for object recognition and optical flow. These are two important application capabilities for future embedded vision environments such as on-board unmanned aerial vehicle (UAV) video processing.

For object recognition, we identified convolutional neural networks as a target for analysis, implementation and optimization. We measured an energy efficiency (GFLOPS/W) of 0.70 for a convolutional neural network with five convolutional layers performing an object recognition task. This is 21% of the theoretical compute bound for this Haswell processor.

We also identified three important optical flow approaches: Horn-Schunck, Lucas-Kanade, and Brox. We analyzed and implemented optimized versions of these approaches. They achieve 0.0338, 0.0103 and 0.0203 GFLOPS/W respectively. We achieve 7.9% of the theoretical bandwidth bound, assuming no cross-iteration memory optimization, for Horn-Schunck optical flow using the Jacobi solver, and 9.8% of the bandwidth bound for the conjugate-gradient solver. To improve performance on optical flow, we will focus on increasing bandwidth utilization. We were not surprised by the low performance because that the code was previously tuned for GPUs. Also, we expect the Intel OpenCL compiler will generate better results than the AMD compiler, which is our current setup. After improving memory bandwidth utilization, we plan to add cross-iteration memory optimizations such as blocking and tiling the Jacobi solver and communication-avoiding linear solvers.

Finally, we identified a meaningful accuracy vs. runtime tradeoff between the three optical flow approaches. Each method shows distinct advantages in terms of the accuracy-runtime tradeoff compared to the other two. We conclude from this that we should continue to develop all three optical flow methods in the future.

A Horn-Schunck Derivation

The following is a derivation of Horn-Schunck optical flow method. This derivation can be found in the original Horn-Schunck paper [3], or numerous other introductions to optical flow [3].

Optical Flow Equation: Assume that all pixels in an image are displaced from time $t$ to time $t + \Delta t$ by $\Delta x$ and $\Delta y$. This is known as the gradient constancy assumption.

$$I(x, y, t) = I(x + \Delta x, y + \Delta y, t + \Delta t)$$ (7)
Next, approximate \( I(x + \Delta x, y + \Delta y, t + \Delta t) \) using a Taylor series expansion.

\[
I(x + \Delta x, y + \Delta y, t + \Delta t) \approx I(x, y, t) + \frac{\partial I}{\partial x} \Delta x + \frac{\partial I}{\partial y} \Delta y + \frac{\partial I}{\partial t} \Delta t
\]

\[
I(x + \Delta x, y + \Delta y, t + \Delta t) \approx I(x, y, t) + \frac{\partial I}{\partial x} \Delta x + \frac{\partial I}{\partial y} \Delta y + \frac{\partial I}{\partial t} \Delta t + \frac{\partial I}{\partial x} \Delta y \Delta t + \frac{\partial I}{\partial y} \Delta x \Delta t + \frac{\partial I}{\partial t} \Delta t \Delta t
\]

This leads to the optical flow equation. The vectors \( u \) and \( v \) are the \( x \) and \( y \) pixel displacements respectively.

\[
\frac{\partial I}{\partial x} u + \frac{\partial I}{\partial y} v + \frac{\partial I}{\partial t} = 0
\]

Written again in terms of the gradient images \( I_x, I_y, I_t \).

\[
I_x u + I_y v = -I_t
\]

**Smoothness constraint:** Notice that the optical flow equation is under-determined. There is one equation for each pixel and two unknowns. Additional constraints need to be introduced in order to solve the aperture problem. The Horn-Schunck method introduces a smoothness constraint, penalizing changes in the flow vectors across the image. The formulation is an energy functional that we seek to minimize.

\[
E = \int \int (I_x u + I_y v + I_t)^2 + \alpha^2 (\| \nabla u \|^2 + \| \nabla v \|^2) dxdy
\]

Since \( u(x, y) \) and \( v(x, y) \) are functions, \( E \) is a functional. We use the Euler-Lagrange equations to find fixed points of this functional. The variables \( u_x, u_y, v_x \) and \( v_y \) are the partial derivatives of the flow vectors. \( F \) is the integrand of \( E \).

\[
\frac{\partial F}{\partial u} - \frac{\partial}{\partial x} \frac{\partial F}{\partial u_x} - \frac{\partial}{\partial y} \frac{\partial F}{\partial u_y} = 0
\]

\[
\frac{\partial F}{\partial v} - \frac{\partial}{\partial x} \frac{\partial F}{\partial v_x} - \frac{\partial}{\partial y} \frac{\partial F}{\partial v_y} = 0
\]

These are the terms we need to plug in to the Euler Lagrange equations for Horn-Schunck optical flow.

\[
\frac{\partial F}{\partial u} = 2I_x(I_x u + I_y v + I_t)
\]

\[
\frac{\partial F}{\partial v} = 2I_y(I_x u + I_y v + I_t)
\]

\[
\frac{\partial F}{\partial u_x} = 2\alpha^2 u_x
\]

\[
\frac{\partial F}{\partial u_y} = 2\alpha^2 u_y
\]

\[
\frac{\partial F}{\partial v_x} = 2\alpha^2 v_x
\]

\[
\frac{\partial F}{\partial v_y} = 2\alpha^2 v_y
\]
\[
\frac{\partial}{\partial x} \frac{\partial F}{\partial u} = 2\alpha^2 u_{xx} \tag{21}
\]
\[
\frac{\partial}{\partial y} \frac{\partial F}{\partial u} = 2\alpha^2 u_{yy} \tag{22}
\]
\[
\frac{\partial}{\partial x} \frac{\partial F}{\partial v} = 2\alpha^2 v_{xx} \tag{23}
\]
\[
\frac{\partial}{\partial y} \frac{\partial F}{\partial v} = 2\alpha^2 v_{yy} \tag{24}
\]

After plugging in these terms, these are the Euler Lagrange equations for Horn Schunck optical flow.

\[
I_x(I_xu + I_yv + It) - \alpha^2(u_{xx} + u_{yy}) = 0 \tag{25}
\]
\[
I_y(I_xu + I_yv + It) - \alpha^2(v_{xx} + v_{yy}) = 0 \tag{26}
\]

We discretize the image derivatives \( I_x \) and \( I_y \) using a \([1/12, -8/12, 0, 8/12, -1/12]\) filter. We discretize the image derivative \( I_t \) and the partial derivatives of the flow vectors \( u \) and \( v \) using a \([-1,1]\) filter. The term \((u_{xx} + u_{yy})\) becomes a familiar 5-point stencil as shown in the following algebra.

\[
u_x = u(x+1,y) - u(x,y) \tag{27}
\]
\[
u_{xx} = (u(x+1,y) - u(x,y)) - (u(x,y) - u(x-1,y)) \tag{28}
\]
\[
u_{xx} = u(x-1,y) - 2u(x,y) + u(x+1,y) \tag{29}
\]
\[
u_{xx} + \nu_{yy} = u(x-1,y) + u(x,y-1) + -4u(x,y) + u(x+1,y) + u(x,y+1) \tag{30}
\]

References


